

www.DataSheet4U.com

CMOS GATE ARRAYS HD61 SERIES DESIGNER'S MANUAL AND PRODUCT SPECIFICATION

#U80



B

LIFE SUPPORT POLICY

Hitachi's products are not authorized for use as critical components in life support devices or systems without the prior express written approval of Hitachi America Ltd.

- 1. Life support devices or systems are devices or systems which,
 - (A) are intended for surgical implant into the body, or
 - (B) support or sustain life, and whose failure to perform, when properly used in accordance with instructions for use provided in the labeling, could reasonably be expected to result in a significant injury to the user.
- 2. A critical component is any component of a life support device or system whose failure to perform can be reasonably expected to cause the failure of the life support device or system, or to affect its safety or effectiveness.

When using this manual, the reader should keep the following in mind:

- 1. This manual may, wholly or partially, be subject to change without notice.
- All rights reserved: No one is permitted to reproduce or duplicate, in any form, the whole or part of this manual without Hitachi's permission.
- Hitachi will not be responsible for any damage to the user that may result from accidents or any other reasons during operation of his unit according to this manual.
- 4. This manual neither ensures the enforcement of any industrial properties or other rights, nor sanctions the enforcement right thereof.
- 5. Circuitry and other examples described herein are meant merely to indicate characteristics and performance of Hitachi semiconductorapplied products. Hitachi assumes no responsibility for any patent infringements or other problems resulting from applications based on the examples described herein.

September 1985

[©]Copyright 1985, Hitachi America Ltd.

Printed in U.S.A.

PAGE

SECTION I

www.DataSheet4U.com

HD61 SERIES DESIGNER'S MANUAL	
1. LOGIC DESIGN	1
2. DELAY TIMES	27
3. TEST VECTORS	33
4. GATE ARRAY FORMATTED LOGIC DRAWING	39
5. GENERATING TEST VECTORS	49
6. HITACHI DOCUMENTATION	59
SECTION II	
HD61 SERIES PRODUCT SPECIFICATION	71

 Ω

INTRODUCTION

This document is provided as guidance to design a system using a HITACHI CMOS GATE ARRAY. The reader is advised to use both sections of this document in order to gain a full understanding of the techniques and procedures involved.

E



www.DataSheet4U.com

1. LOGIC DESIGN

T+Y	Gate Utilization factor	2
1.2	Delay times	3
1.3	Fan-out limitations	4
1.4	Simultaneous transition of output signals	5.
1.5	Non-connected pins	6
1.6	Wired 'OR' restrictions	7
1.7	Restrictions on linked input connections	8
1.8	Chopper circuit	9
1.9	Latch built with discrete logic gates	10
1.10	Floating three-state bus	11
1.11	Initialization	12
1.12	Split grouping of cascaded counters	13
1.13	Test clock	14
1.14	Test pins	15
1.15	Test circuit	16
1.16	Drivability of pre-buffers	1,7
1.17	Construction of three-state buffer	18
1.18	Hazardous conditions	19
1.19	Operational restrictions on Flip-Flop and Latch	20
1.20	Serial synchronous data transmission	21
1.21	Differences in operation between silicon and simulation	22
1.22	Pin assignment	23
1.23	Miscellaneous	24

The usual arrangement for laying out gates in a GA is to organize the gates in rows of cells, with gaps between the rows of comparable thickness to the cells so that sufficient space is allowed for connections to be made between cells. Since layout is performed by a CAD system automatically, as the design becomes more complex and the total gate count of the logic design approaches the total number of gates available, routing of the interconnect becomes immensely difficult, requiring a long engineering period.
To guard against this situation, HITACHI recommends that the gate count utilization is limited to no greater than 90% for 61J, K, L and MM series (without RAM) and 80% for 61MM series with RAM.

Master	Gates	Max. Usage
HD61J HD61K HD61L	504 1080 1584	450 970 1420
HD61MM without RAM HD61MM	2496 2496	2150 2000
with RAM		-

EX.1 Calculation of equivalent gates

Macrocell	Equiv. gates (G)	Number used (N)	Gate Count (G × N)
FD	6	4	24
NA1	0.5	4	2
NA2	1	9	9
NA3	1.5	10	15
NA4	2	1	2
NA6	4.5	5	22.5
EOR	2.5	4	10
Total			84.5

Because input and output buffers are fixed at the peripherals of the gate array, they are ignored in the calculation of total gate count.

EX.2 TTL logic based diagram

When gate count is calculated with an equivalent TTL circuit, refer to the conversion table in section 6.8.

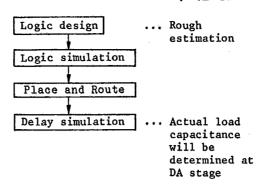
However, care should be taken when interpreting this table since in some situations the gate count assumes that the whole of the TTL function is utilized, whereas in reality only 70-80% might be used.

At the time of logic design approximate calculations should be made of the delays that are likely to occur through various logic paths in the GA.

Logic paths which exhibit delays likely to cause problems to the customer's design should then be identified to HITACHI as 'Critical Paths'.

Following simulation, auto routing will yield actual delay times that will occur in the GA which can then be checked by the customer to determine that the design is correct.

The procedure for Delay specification is as follows:



(1) Rough estimation

During logic design, estimate the total delay times along the critical path.

td = to + K × C_L $C_L = 0.4 \times EFO$ $EFO = \sum_{F,0}$ When calculating, refer to the Macrocell library for delay parameters. Those figures are typical, i.e. they represent typical loading conditions for average routing length and fan out.

Once the logic is autorouted, more precise delay times for the user defined paths are available from the DA for further analysis.

(2) Designation of critical paths

A critical path should be designated when the logic diagram is released to Hitachi.

Delay check report is derived from the actual routing information.

(3) Evaluation of prototype devices

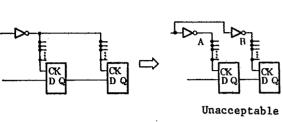
Hitachi debugs the prototypes using the test vectors supplied by the user. The user then evaluates these in the actual system in order to check that timing margins are within the design specification.

If delay times meet the requirements of the design, then no limitations are placed on fan out. However, in the case of a clock driver for F/F, when the fan out increases tr/tf will degrade until eventually a malfunction occurs.

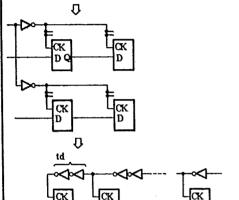
Maximum fan out for clock drivers

Power inverter 20 Others 10

< Example> Clock skew



When a large fan out is required, it is necessary to separate the loads by providing separate drivers. However, in the example shown, the driver may cause a malfunction due to the clock skew between A and B.



Group F/Fs which share the Data in serial.

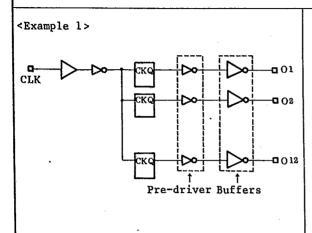
Delay the clock supplied to the front-end F/F rather than the one which follows. $td \ge 5$ ns

14 SIMULTANEOUS TRANSITION OF OUTPUT SIGNALS

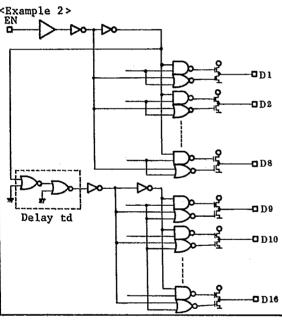
T-42-11-09

When many outputs change simultaneously, high transition currents are drawn through the ground line. This causes the ground level to swing due to the inductance existing in the chip. Consequently this leads to a reduction in the $V_{\rm IH}$ margin. As a general rule, for the purposes of reducing ground rise, Vcc and ground lines in PCB should be ept wide and short. And simultaneous transitions on output signals should be restricted according to the values shown in the table.

Pre-driver	Max. simultaneous transition
Inverter/NAND	12
2 input NOR	16
3 input NOR	18



In example 1, inverter pre-drivers are used. In this case, the output should be limited to a maximum of 12.



This example contains 2 sets of 8-bit bus drivers. Dummy gates are inserted to delay one set of bus drivers from the other.

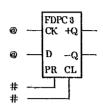
In this case, data should be consistent during the time the buffer is enabled.

Note: Test vectors should also be stable during this period.

td \geq 10 ns. Typ.

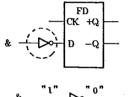
The input pins of unconnected macrocells are automatically modified by the DA and connected to either Vcc or OV. These levels are determined by DA program according to the rules described in the cell library. (Refer to Macrocell Library). Inputs which are not described in the Macrocell library will be connected to "1" (Vcc level) and "0" (GND level) respectively for AND/NAND and OR/NOR Macro. Inputs to other macros must be defined by the user to either "1" or a "0". Under no circumstances must these be left floating.

<Example 1>



All unconnected pins will be modified as follows:

<Example 2>



The DA automatically connects the D input of macrocell FD to "1" level. When an "0" level is required, an inverter should be added i.e. as shown in example 2.

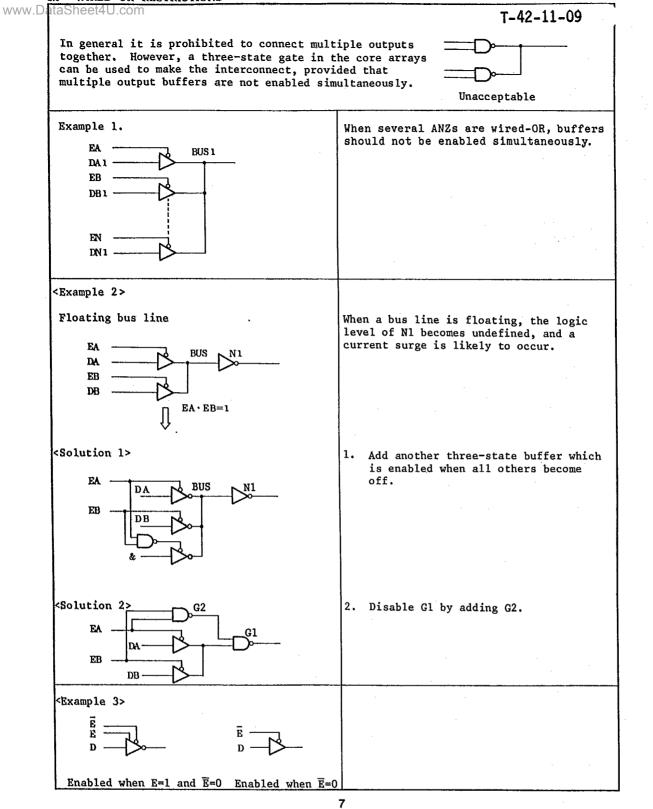
<Example 3>

Example of macrocells which are not described in macrocell library.

AND/NAND ... "1"

OR/NOR ... "0"

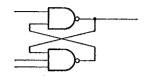
EOR/ENOR ... "O"



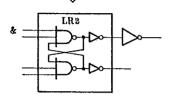
Unacceptable

Latch circuit or gate circuit with feed back loop should not be constructed with NAND or NOR logic gates. Instead, flip-flops or latch macrocells should be used.

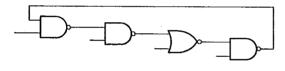
<Example 1> Latch



Should be replaced by



<Example 2> Looped circuit with
primitive logic gates

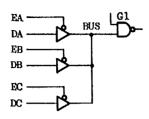


Unacceptable

When internal bus could take three-state, a surge current may flow in the gate which is driven by this bus line because of the intermediate input level. In order to prevent these unacceptable conditions the following is suggested:

- (1) Do not float the bus line
- (2) Gate the successive gates by synchronizing with an enable signal of the bus line

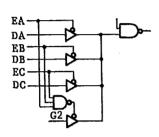
<Example 1>



Enabled when E=0

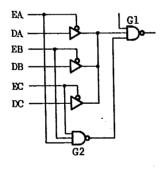
When the condition $E_A=E_B=E_C$ occurs a transitional surge current is likely to occur in successive G1.

Alternative 1



Addition of a dummy NAND gate, G2, forces the extra three-state buffer to "O" level, when E_A , E_B and E_C disable the bus line.

Alternative 2

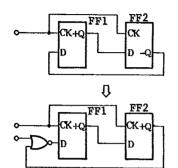


The additional gate will disable Gl when the bus enters the high-impedance state.

Just after power is applied, the logic states of latches and F/Fs are undefined. Logic should then be initialized with the appropriate test vectors, otherwise the design is not testable. During logic simulation, logic states are also undefined at t=0. As a consequence, the gates in the core have to be initiated step by step, starting with the gates nearest to the GA inputs until finally all gates are in a known state.

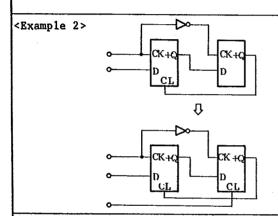
<Example 1>

Example of a circuit which cannot be initialized



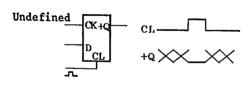
When an output of FF2 is fed back to the front FF1, it is not possible for these FFs to be initialized.

By providing an external control to the data input of FF1, FF1 can be initialized on the first clock and FF2 on the second clock.



In the actual device, this circuit will be initialized when the data input is "O" and a clock occurs. However, for the simulation a problem will occur because the state of F/F is undefined when either CK, CL or PR are undefined.

<Example 3>

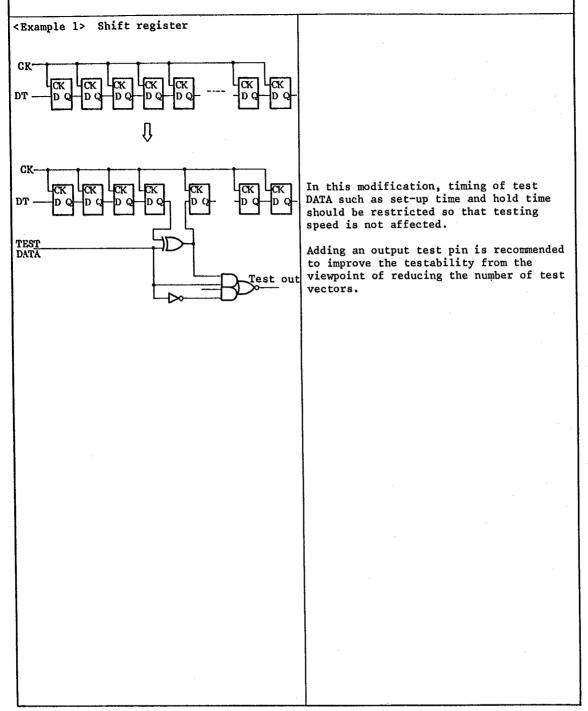


CK, CL & PR should be defined states during initialization.

Indefinite X Indefinite

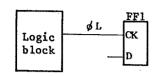
Even though the relation between X and Y is evident, both X and Y need to be a "1" or an "0".

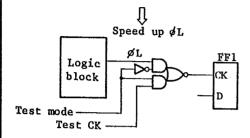
Because the hardware of the tester restricts the number of test vector steps, it is suggested that in the case of many serially cascaded counters, these counters should be split into groups and controlled from an external test pin.



In order to minimize test vectors, (a situation that occurs when the logic configuration of a particular logic block results in long periods, relative to the maximum clock rate, between logic state changes), an effective method that can be used is to include in the logic of the GA a test circuit, controlled (enabled) from an external pin.







In the example, the repetition rate of ϕL is much slower compared to maximum clock rate. To overcome this problem, a multiplexer can be placed between the logic block and FF1.

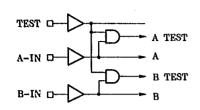
If the multiplexer switches when ϕL and TEST CK are both set to a logic "1", a glitch is likely to occur at the CK input to FF1.

Consequently, it is necessary to switch the multiplexer when ϕL and TEST clock are both set to '0' levels.

Note: Function of the multiplexer should not degrade the margin for set-up and hold times of FF1. In the event that insufficient pins are available for including test pins, it is suggested that normal function input and output pins would be shared with those used for test functions.

<Example 1>

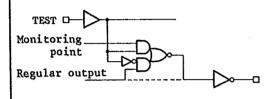
Combining test inputs with normal inputs:



"A TEST" and "B TEST" are enabled when "TEST" goes to '1'. In this example regular functional pins A-IN and B-IN are used as test pins in conjunction with "TEST" pin.

<Example 2>

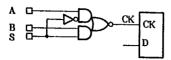
Sharing test outputs:

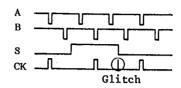


When "TEST" swings to "1" level certain core gates can be monitored. This can be very helpful in shortening the test vectors required, especially when the technique of monitoring the innermost core logic directly at the output pin via a multiplexer is applied.

A test circuit is often implemented in order to improve the testability of the logic. The implementation necessarily requires that the timing of the design is such that glitches which can occur at the switching of modes are prevented, and that the margin of set-up and hold times for the F/F is met.

<Example 1>



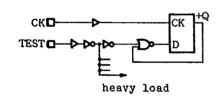


Switching the mode "S" during A = B = "1" causes a glitch due to the delay at inverter. Consequently, switching at A=B=0 is required.

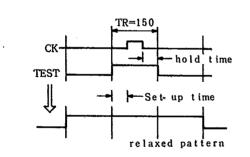
Note: Glitches should be prevented from occuring in the simulation even though they are unlikely to occur in the actual device.

Œxample 2>

Margin for set-up time.



F/F may malfunction due to the shortage of set-up time.

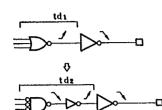


Insertion of dummy step may provide a solution.

T-42-11-09

Since the gate capacitance of an output buffer is relatively large, the drivability of pre-buffer should be considered when designed in. If the timing specification is tight, a NAND gate or an Inverter is recommended as a pre-buffer.

Example 1>



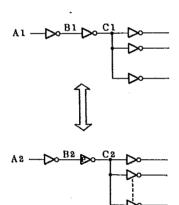
td1 = 16.5 ns

td2 = 9.5 ns

This calculation shows that a NAND gate is faster than a NOR gate.

(Realization of this modification is necessary to invert the related inputs).

Example 2>

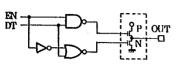


Power gates are effective for minimizing the delay time if employed in a heavy loaded path. If, however a path is lightly loaded, delay times can be longer than expected as the load factor of power gates is larger than that of an inverter.

A three-state buffer cell does not provide the control gate but only a pair of C-MOS transistors. Control gates should be prepared from core gates.

An I/O common buffer needs to be designed in the same way since control gates are not included either.

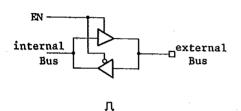




Output buffer cell

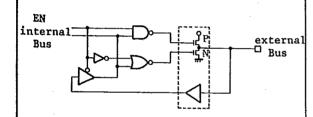
In the example the output buffer is enabled when EN = "1" and the output data has the same phase as an internal DT signal.

<Example 2>

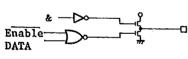


External bus line should not be floating since the input buffer in an I/O common buffer cell is always enabled.

When preparing the test vectors, I/O mode switching timing should be given careful consideration. (Refer to 5.2 page 49)



<Example 3>



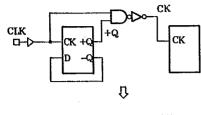
Prohibited

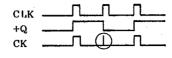
Use of a three-state buffer as an open drain buffer with a built in pull up resistor is prohibited.

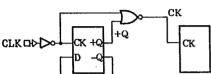
Timings relating to F/F logic should be carefully designed so that pulse width, set-up time and hold time are kept within adequate margins. (The test vectors for F/F types should be given careful consideration too).

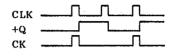
<Example 1>

Glitch on clock:



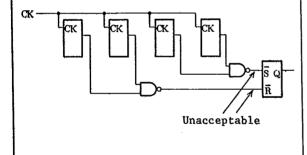






<Example 2>

Synchronous counter:



As outputs from synchronous counters are not exactly synchronized due to the routing path variations and differences in the processing parameters during wafer fabrication, logic should be designed so that such variations have no effect.

If a simultaneous transition occurs on multiple inputs, glitches are easily predictable. However, these are not always observable in the logic simulation.

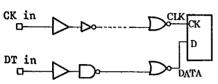
In the logic design cycle, these logic dependent glitches and critical timing should be studied thoroughly.

All F/F and latch types require that the following list of parameters be satisfied. The figures listed are the data referenced at the macrocells excluding the delay times along the path.

(Refer to 1.20 regarding synchronous data transmission).

Parameters	Symbol	Times (ns)			
Clock pulse width	twck	50 min.			
Data set up	tsu	40 min.			
Hold time	th	20 min.			
CLR, PR to clock	tr	50 min.			
Pulse width CL, R	twCL	50 min.			
PR, S	twPR	50 min.			
R to S	tRS	50 min.			
S to R	tSR	50 min.			

<Example 1>

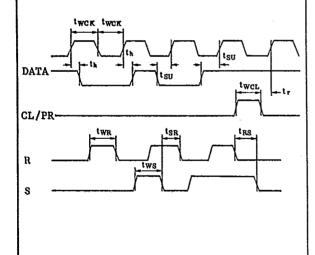


Delay times from CK in to CLK and DT in to DATA should be included when the timings of "CLK" and "DATA" are analyzed.

<Example 2>

It can be seen from the list of timing restrictions that F/F and latch types are allowed to operate up to approx. 10 MHz. However, it may be feasible to handle up to 15 MHz depending on the structure of the logic and the specification.

Timing chart:

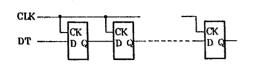


_

An exceptional application of a F/F type is when used as a shift register or synchronous counter supplied with a single clock and used in serial synchronous data transmission mode.

<Example 1>

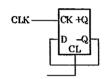
Shift register:



The clock can be shared for multiple F/Fs if serially cascaded. However, splitting the clock drivers may cause a malfunction of F/F types because of clock skew.

<Example 2>

Counter:

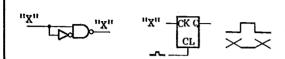


This design is acceptable despite a short hold time which is less than the specification.

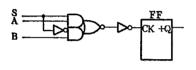
Reasons for discrepancies between the simulation and actual silicon are as follows:

- Recause of limitations in the simulator model some gates cannot be initialized. (Ref. 1.10)
- (2) Glitches and spikes while possible in a simulation, do not necessarily occur in the actual device.
- (3) Glitches can occur in the silicon due to layout or process variation.





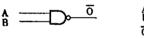
<Example 2>



In the simulation a glitch is observed at the switching during A=B=1, and it will enable F/F.

In the silicon, F/F may not respond to a narrow glitch. Consequently, it is prohibited to prepare test vectors that are conditional on the glitch to enabling successive gates.

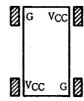
<Example 3>





Even though the glitch is not observed in simulation, since there is no process variation and the timing is exact, in the actual device these variations may occur and so too the resultant glitches. All I/O pins can be assigned by the user except power supply pins.

It is recommended that any noise sensitive functions such as clock for F/F and latch and asynchronous clear inputs should be assigned close to the ground pin. On the contrary, pins which are a source of noise, such as 8-bit bus drivers which switch simultaneously, should be tracked close to the other ground pin.



<Example 1>

It is recommended that pins that are not connected (NC) should not be utilized, but should be connected to power supply pins or be left open.

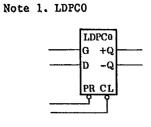
In Gate Arrays, since all processing of manufacturing is standardized, NC pins are also bonded inside the mould. Consequently, excessive voltage at the package terminals may cause malfunction or damage the chips.

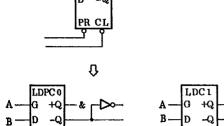
RESTRICTION ON PIN ASSIGNMENT

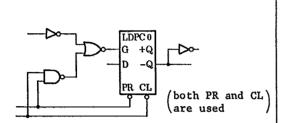
Assignment of power supply pins is fixed. Modification from this standard should be consulted of its feasibility for each package respectively.

		DP-28	DP-40	DP-42	DP-64	FP-54	FP-80	FP-100	DC-28	DC-40	PGA-72
HD61J	GND	1,15	1,21	1,22		1,28			1,15	1,21	
	VCC	14,28	20,40	21,42		27,54			14,28	20,40	
HD61K	GND	1,15	1,21	1,22	1,33		33,74		1,15	1,21	
	VCC	14,28	20,40	21,42	32,64		31,72		14,28	20,40	
	NC						1,3,22,24,41,43,62,64				
HD61L	GND	1,15	1,21	1,22	1,33	Γ	34,74		1,15	1,21	44,72
	VCC	14,28	20,40	21,42	32,64		32,72		14,28	20,40	60,65
	NC						2,3,22,23,42,43,62,63				
HD6IMM	GND		1,21		1,33			29,41,79,91	1,15	1,21	
	VCC		20,40		32,64			2,39,52,89	14,28	20,40	

Vcc: 5V ± 5% GND: O.V





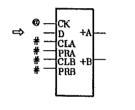


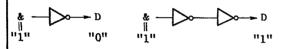
(PR is fixed)

- While clear is enabled, if D = 1 and G goes high both +Q and -Q become "O".
- While preset is enabled, if
 D = "0" and G goes high both
 +Q and -Q become "0".

Note 2. ZSRCP3

(CL is fixed)

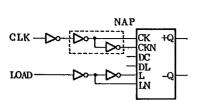




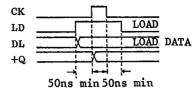
D is requested to be fixed "0".

D is requested to be fixed "1".

- D input of macorcell ZSRCP3 is not automodified when nonconnected.
- The preset function is enabled synchronously to clock, while clear function is enabled asynchronously.



- CK and CKN need to be driven by a power inverter cell, NAP.
- 2. Timing of loading data.

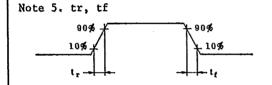


FDL operates the same way as ordinary F/F under the condition of L=0 and LN=1.

Note 4. Simultaneous enabling of clear, preset, set or reset

Cell	PR(S)	CL(R)	+Q	-Q
LRS0	0	0	0	0
LRS3	1	1	1	1
LR2S20	0	0	0	0
LR2S23	1	1	1	1
LDPC0	0	0	0	0
FDPC3	1	1	1	1
FJC1	0	0	0	0
ZSRCP3	1	1	0	- .

Disabling these signals simultaneously is prohibited because the logic state will be undefined after the signals are released.

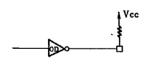


The following waveform is recommended:

 $tr \le 100 \text{ ns}$ $tf \le 100 \text{ ns}$

Note 6.

www.DataShee34UF62m



An external pull-up resistor applied to an open-drain buffer should be biased within the same range of Vcc or GND as is supplied to the gate array chip.

26

2. DELAY TIMES

T-42-11-09

2.1	Calculation of delay times	2
2.2	Calculation of loading capacitance	2
2.3	Variation of delay times	3
2,4	Example of calculation	3

Precise delay times for all logic paths in the GA will ultimately become known once the DA has auto-routed the interconnect. However, before this stage is reached the designer should check, by hand calculation, that the delay times are realizable.

(Refer to the macrocell library for the values "to" & "K".)

Formula

$$\begin{vmatrix} t_{pLH} = t_{oLH} + K_{LH} \cdot C_L \\ t_{pHL} = t_{oHL} + K_{HL} \cdot C_L \end{vmatrix}$$

LH: Low to High HL: High to Low

<Example 1> FDPC3

How to calculate delay

	Symbol No.	Delay					
Symbol	· 유 ·	Input	Out- put	tpln (ns)		tpHL	(ns)
	Sy	Name	Name	to LH	K _{LH}	to HL	K _{HL}
	c	СК	+Q	4.6	2. 4	3. 9	1. 6
	3	CK	- Q	5. 2	2. 2	5. 5	1. 4
FDL	В 4	CK/ CKN	+Q -Q	2.5	2. 5	3. 1	2.1
CKN +Q-		L/LN		7.7		9.0	
□ IN −Q □	D 4	CK/ CKN		4.3		3. 6	
.		L/LN		10.7		8. 5	
(СK		5. 4		4, 2	1. 7
FDPC3 		CL	+ Q	2. 5	4.4	1. 2	1. 7
p -Q-	С	PR		5, 5		_	_
PR CL	"	СК		6.8		6. 8	1.5
	ا استنها	CL	- Q	5. 8	4.4	_	_
		P R	<u> </u>	4.0		1.5	1.5

tLH: Low to high at(F + Q) when calculating from a rising CK edge.

- 1. Select the related inputs (CK, +Q)
- 2. Pick out the figures $(t_{OLH} = 5.4, K_{LH} = 4.4)$
- 3. Calculate according to the formula t_{pLH} = 5,4 + 4,4 × C_L (ns)

Loading capacitance is obtained from the summation of input capacitance and stray capacitance which exists in the inter-connection metal layer. The DA supplies the ultimate precise value of capacitance, however, during the logic design phase capacitance is estimated as 0.4 pF/inverter gate.

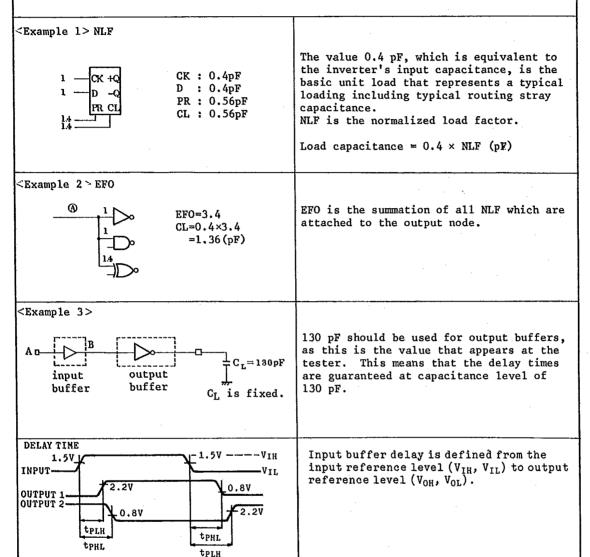
Effective fan-out

 $EFO = \sum_{F,O} NLF$

Total capacitance

 $C_L = 0.4 \times EFO (pF)$

NLF: Refer to cell library



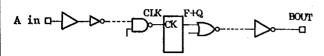
Typical delay times are obtained by the formula shown in 2.1. These values vary depending on the operational conditions such as environmental temperature, power supply tolerance, variation of wafer processing and as a result of errors existing within the calculation.

Variational range:

Min. =
$$0.3 \times typ$$
.

Max. =
$$2.2 \times \text{typ}$$
.

<Example 1> Critical Path:



Indicate the points of origin and termination when the path is unique.

Delay checking list

Designate the interadjacent nodes when the signal has multiple routes between the original and the termination.

```
No Critical path Requirement(ns)

A in √∿ B out \ 180
```

ex: AIN / \CLK / \((F + Q) \\ \\ \Bout \\

<Example 2>

Delay times after auto-layout

Delay times of critical paths which are calculated using actual loading capacitance, are supplied after auto-layout (place & route).

Variation after layout:

Min. =
$$0.4 \times typ$$

$$Max. = 2.0 \times typ$$

The procedure for delay calculation for the logic diagram in section 6.6 is shown below.

Example 1> Path 1:

Clock / ∿ FFA + Q\ ∿ CAA / ∿ A 306 \ ∿ RCA \

						•	
No.	input ∿ output	Cell	(A)	®ĸ	© _{EFO}	$\mathbb{O}^{\mathrm{C}^{\Gamma}}$	td
1	Clock -ClockI /	ANI	7.7	2.0	4	1.6	10.9
2	ClockI /-FFA+Q	FD	3.9	1,6	2.4	0.96	5.4
3	FFA+Q\-CAA	EOR	4.2	4.5	9.0	3.6	20.4
4	CAA/-A306\	NA6	6.0	2.0	1	0.4	6.8
5	A306\-A332 /	NA2	1.3	3.6	4.7	1.88	8.1
6	A332/-RCA	NAOT	4.0	0.34	-	130	48.2
7	total						99.8

Maximum delay time = 99.8×2.2 = 220 (ns) EFO = Σ NLF fan out $C_L = 0.4 \times EFO$ td = to + KC_L

<Example 2 > Path 2:

Clock I / ∿ FFA+Q / ∿ CAA / ∿ A317 \ ∿ FFD+2 /

No.	input ∿ output	Cel1	⊕ ₀	® _K	© _{EFO}	Φ^{Γ}	td
1	ClockI/10FFA+Q/	FD	4.6	2.4	2.4	0.96	6.9
2	FFA+Q/∿CAA/	EOR	4.2	4.5	9	3.6	20.4
3	CAA/∿A317 \	NA6	6.0	2.0	1	0.4	6.8
4	A317\~FFD+2 /	NA3	2.0	3.9	1	0.4	3.6
5	total						37.7

Maximum delay time = 37.7×2.2 = 83 (ns)

32

3. TEST VECTORS

3.1	Definition of test vectors	34
3.2	Fault injection	35
3.3	Fault coverage	36
3.4	Initialization	37
3.5	Timing chart and binary coding	38

www.DataSheet4U.com

A test vector is the expression of translated functions which should be implemented in the GA. Test vectors are utilized not only for the functional verification of the logic but also for final testing during the mass-production phase.

<Example 1> Contents of test vectors:

- 1. DC parameter test
- 2. Function test
- 3. High speed operational test.

Fundamentally, functional test vectors are the only ones which are usually required from the customer. However, a high speed operational test vector may be required if the functional test vectors are not adequate.

- * DC parameter test:
 Output level (V_{OH}, V_{OL}), input leakage,
 (I_{LI}), output leakage of three-state
 buffer and/or open drain buffer (I_{LO})
 are tested. Since these items are
 automatically compiled from the
 function test vector, it should contain
 all logic states within 2000 steps.
- * Functional test:
 For the functional test the logic function fault coverage should be greater than 95%.
 It is suggested that input/output timing is similar to that required for the actual application.
- * High speed operational test:
 If requested, high speed tests will
 be applied and may be abbreviated if
 the required speed is slow. If the
 functional test vectors are not
 applicable to this test, another
 set of test vectors will be requested:
- When there are major differences in delay times from multiple outputs.
- (2) When strobe timing is critical because of high test rate and also because of large delays due to the 130 pF loading capacitance.

<Example 2> Timing test:

- 1. Timing tolerance (set up, hold time)
- 2. Cycle time
- 3. Absolute delay times
- 4. Time difference between output pins
- input timing
- test rate
- → strobe timing
- ← not applicable

The following faults are injected (single stuck at "1" or stuck at "0".

- 1. Input of all logic gates "0" "1"
- 2. Output buffer "0" "1"
 excluding: internal gates of macrocell
 : indistinguishable equivalent
 faults (collapsing)

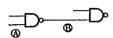
Injections

AND/NAND	"1"
OR/NOR	"0"
EOR	"0" "1"
D input (F/F)	"0" "1"
Output pin	"0" "1"

Definition:

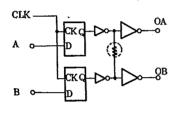
"0" injection = fault stuck at "0"
"1" injection = fault stuck at "1"

<Example 1> Equivalent fault:



The stuck at "0" fault at node (A) is equivalent to a stuck at "1" fault at node (B) In this case stuck at "0" fault at (A) is not injected.

<Example 2>0/1 fault injection:



A fault is detected by comparing the fault injected simulation with a normal (fault free) simulation.

However, faults may go undetected if incorrect test vectors are applied to the input of a particular logic function.

In the left figure, a fault shown will not be detected by the test pattern (A,B) = (0,0), (1,1) even though this gives a coverage of 100%. To detect this fault requires the pattern (A,B) = (1,0), (0,1) to be applied.

Faults are ultimately detected at the output terminals of the GA, after the error in the logic function during the simulation is propagated to the terminals.

Fault coverage = detected faults total injected faults × 100%

Since fault coverage has the dominant influence on the quality of the products, the acceptable coverage level is at least 95%.

Acceptable coverage ≥ 95%

<example 1=""></example>	A 0/1 stuck fault is detectable when the node provides both 0/1 logic levels and the condition can be observed at the related output pin. This concept should be considered when designing test vectors.
<example 2=""></example>	
Transition characteristics:	Transition times (Tr/Tf) are not involved in fault simulation but DC 0/1 levels are considered. Delay times should be tested by a separate functional speed test.
<example 3=""></example>	
Detection of level-fixed node:	Macrocell pins which are clamped to Vcc or OV are excluded from the fault simulation.
<pre><example 4=""> Detection of clear signal of counter:</example></pre>	Detection of signal at the clear pin of an F/F is achieved by the following procedure:
CLK CK CK CL CL	 Set the F/F at "1", first. Enable the clear pins. Supply the clock so that the state of all F/Fs is transferred to the related output pin.
CLK FF1 FF2 FF3 CK CK CK CC C	This example highlights the case which requires a relatively longer test pattern. Also, enabling the clear input after setting all F/Fs to a "l" level is invalid. This is because when the clear input is enabled -Q of FFl goes high. Consequently, it is not possible to distinguish if the F/F2 was activated by CK or the clear signal.

Just after the power is supplied to the GA, the logic states of all the gates are initially unknown. Therefore a requirement of the test vectors is that an initialization sequence should be included as part of the test set.

This is an important consideration, since even though the hardware may allow for external synchronization, the testing method is standardized and cannot be modified.

<Example 1> Number of vector steps:

The synchronous clear function is often employed for counters and for minimization of the gate count.

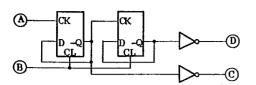
Consequently, F/F types or counters need a well organized test circuit so that the whole circuitry can be cleared by single pulse or within a few steps. Since the number of test vector steps is limited, efficient initialization is absolutely essential.

<Example 2> Splitting of test vectors into separate sets:

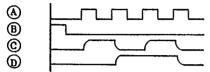
It may be convenient to separate the whole logic into a number of blocks and prepare the test vectors accordingly. However, even in this case, it is recommended that the whole logic is initialized in the first group of vectors that are submitted.

If some portions of the logic remain uninitialized, problems may result when designing the test vectors for another portion. Initialization of the whole logic can help to minimize the steps since the portions which are not being tested may operate anyway.

Test vectors should be binary coded using the formated test vector sheet.

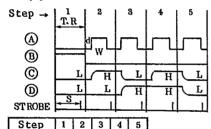


Timing chart



Procedure for converting a timing chart into binary coding.

- 1. Designate the cycle time of the fastest input.
- Specify the clock wave form of (1).
- 3. Define the timings of other signals referring to the clock.
- Check the timing of all outputs assuming zero delay for all gates.
- Calculate the delay times.
 Designate the output which has the largest delay time (or the smallest delay time).
- Decide the strobe timing according to (6).



L						
A	0	1	1	1	1	
В	1	0	0	0	0	
С	L	Н	L	Н	L	
D	L	L	Н	Н	L	Binary coding
		1				or
						, Or
Step	1	2	3	4	5)
Step A	1	2	3	4	5	
Step A B	1 0	2 1 0	3	4	5	
A	1 0 1 L	1	3 L	4 H	5 L	abbreviated

test rate

PP/NP, delay, pulse width PP/NP, delay, pulse width for the clocks & delay time for DT

strobe timing

In order to decide the timing definition, operation of an F/F type should be considered with regard to the followings:

- 1. Set-up time, hold time
- Timing tolerance between the rising edge of the CK and disabling edge of CL or PR inputs on the F/F.

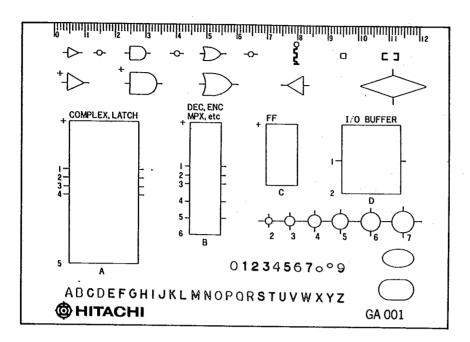
expression

4. GATE ARRAY FORMATTED LOGIC DRAWING

4.1	Drawing sheet and template	41
4.2	Logic symbols	41
4.3	Allocation of names	42
4.4	U-D Macrocell	4:
	Placement of symbols	
4.6	Miscellaneous	4.
4.7	Example of Logic Description	4

The drawing sheet and templates will be supplied by Hitachi to the customer. Size of sheet is A3.

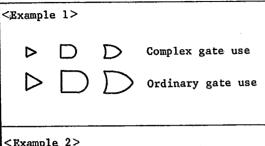
- (1) Refer to the drawing sheet in section 6.4
- (2) Use the template shown below



(Actual size)

It is required that logic diagrams be drawn with HITACHI templates.

- * Refer to the Macrocell library for symbols. Draw the schematic using the same function names, terminal names and the sizes as given in the library.
- Complex gates enclosed with dotted lines can be described in an abbreviated form with only the function name. The input location will be the same as in the non-abbreviated case.

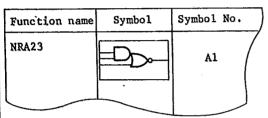


For complex gates enclosed with a dotted line. (Refer to Macrocel'l Library.)

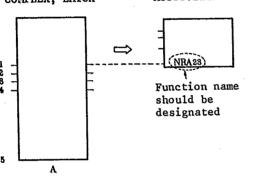
Function name and terminal name are not needed for AND/OR symbols.



NRA23



Abbreviated case COMPLEX, LATCH



The symbol no. in the macrocell library indicates which shape should be adopted while drawing.

An '&' mark should be attached to the level fixed input pins of a macro. They will be then clamped to the appropriate level automatically. (@ : "H" # : "L")

The location of I/O pins of a macro should be approximately identical to that shown in macrocell library.

Each gate should be named using alphabet and signs of +,-. All names should be begun with alphabetical characters. Each line has its own name if the line is independent of another. (i.e. if it is not connected).

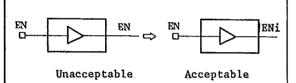
Limitation on naming:

- * Terminals of package ... 2-5 letters
- * Macrocell 2-5
- * Memory macrocell up to 2
- * Three-state gate up to 4

CLK FD FF1+Q CK +Q FF1-Q .

Each macrocell should be named first. Then,

- Output signal would be named with macrocell name itself if macro provides only one output terminal.
- (2) In case of the macros which provide multiple outputs, output signals would be named with the combination of macrocell name and the outputsuffix such as +Q, -Q, +Y, -Y and so on.
- (3) It is allowed to omit naming the line which is terminated on a sheet of diagram.



Separate node should provide unique name.

<Example 3>

EA

DA

EB

DB

Wire-ored lines should not be named respectively but be sharing the name as shown in example.

Definition

<Example 2>

- * Macro function name
- * Macro terminal name
- * Macro name
- * Terminal name

- ...function of macro (ex: NA, NR, FD)
- ...pin function of macro (ex: CK, D, +Q,
- ...unique name given to each macrocell
 ...name given to each node

42

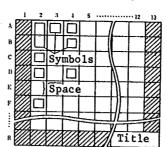
Name Allocation:

- * All nodes inside macro ... within 4 letters beginning with @+/@-
- * Input nodes of macro ... names are constrained to 2 letters of alphabet.
- * Function name ...
 these are constrained to 5 letters
 with the first letter of
 alphabet.
- * Output nodes out of macro ... exclude "@".
- * Symbol size ...
 Width: A size of template
 Hight: (N+1)×2 divisions of the
 drawing sheet, where, N is the
 number of either inputs or
 outputs whichever is greater.

<Example 1> In this case: It is prohibited to define a UD macro for a F/F type or a latch. It is impossible to guarantee the functional operation and the delay Nested UD macros are not supported. <Example 2> I/O buffer cells cannot be included Not applicable in UD macro. **UD** Macro Macro <Example 3> Employ unique names for the nodes in a UD macro. For fixed output names UD Macro of macrocell, such as +Q/-Q of F/F types, the names should be attached like the example showr. Outputs of UD macros should be named with a maximum of 4 letters. <Example 4> The maximum number of I/O connections of either input or output is 40.

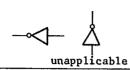
Symbols need to be positioned in point to the right so that signals flow from left to right. Placement in the shaded area is prohibited as in the overlapping of symbols.

In order to ensure sufficient space for routing at least one row in five should be kept blank. Spacing on columns should be for every other column.









Rotation of symbols is not supported.

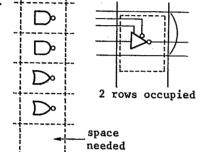
<Example 2>



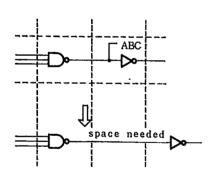
1 column placeable 13th column prohibited

Cell cannot be placed on first and 13th columns.

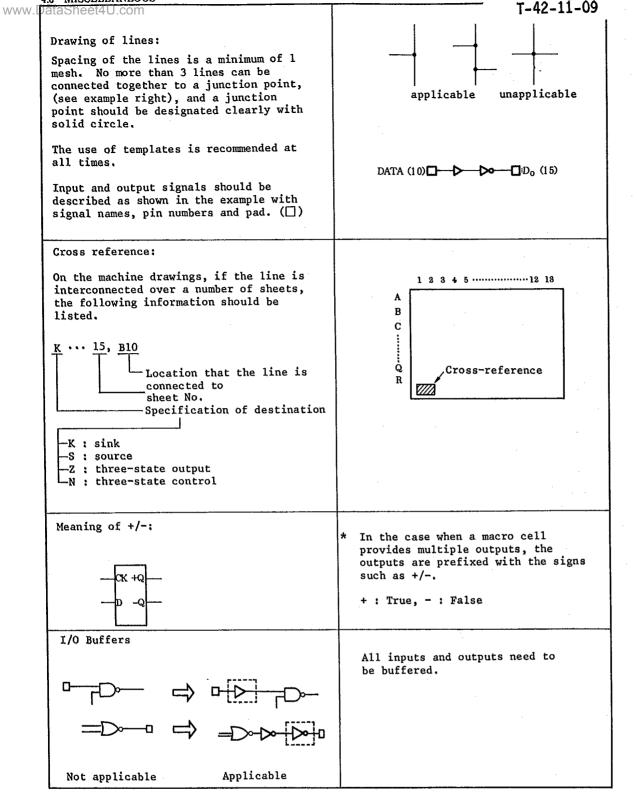
<Example 3>



<Example 4>

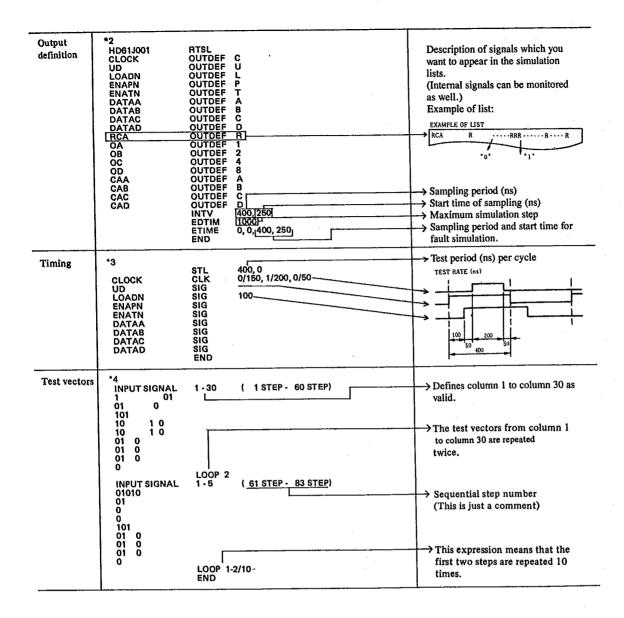


Macrocells should be placed every other column.



■ 4.7 Example of Logic Description

hematic *1 scription HD61J00 HD61J	TRANS	DILP28	Device type, master type package
NB1 @+1 @+2 @	0108 BLOCK 01Q10 NA 01Q11 NA 01Q12 FD BEND	V,W,X,Y,Z @+Q,W,X @+1,Y,Z V,@+2 @+Q,@-Q	UD-macro definition
UP GLOCK DATAA DATAB DATAC DATAD ENAPN	02 INPUT 03 INPUT 04 INPUT 05 INPUT 06 INPUT 07 INPUT 08 INPUT		Description of I/O terminal including pin assignment.
LOADN ENATN OD OC OB OA RCA	09 INPUT	OD1 OC1 OB1 OA1 RCA1	→ LOAND is an input pin located at pin S of the package.
CLOCK1 UDP E021 J1 J2 J3 J4 J5 S6 OA1 OB1 OC1 OD1 RCA1 UDN LOADI ENATI B0701 DAALO DABLO COBO DABLO FFFO1 FFFD1 CAB CAB CAB CAB CAB CAB CAB CAB CAB CA	0182 ANI 0162 ANI 0163 NAOT 01612 NAOT 01612 NAOT 01612 NAOT 01613 NA 0163 NA 0164 NA 0165 NA 0167 NA 0167 NA 0168 NA 0168 NA 0168 NA 0168 NA 0168 NA 0168 NA 0169 NA 0169 NA 0160 NA 0160 NA 0161 NA 0161 NA 0162 NA 0163 NA 0164 NB 0165 NA 0165 NA 0165 NA 0166 NA 0167 NA 0168 NA 0168 NA 0169 NA	CLOCK UD LOADN ENAPN ENATN DATAA DATAAB DATAC DATAC DATAC DATAC FFA-Q FFB-Q FFD-Q P0701 UDP E0201 J2 ENLDN J3, LOADI J4, LOADI J5, LOADI J6, LOADI UDP, CAD ENLDN, CAA ENLDN, CAA ENLDN, CAA ENLDN, CAA ENLDN, CAA ENLDN, CAA, UDN CAD, CAC, CAB, CAA, UDN CAD, CAC, CAB, CAA, UDN CAD, CAC, CAB, CAA, FFD-Q ENLDN, JOSOI, JOSOI, JOSOI, FFB-Q ENLDN, CAC, CAB, CAA, FFD-Q FEA+Q, UDN FFB+Q, UDN FFB+Q, UDN UDN, FFD+Q CLOCKI, B0701, E0201, C0801, DAALO CLOCKI, E0201, E0701, F0801, DABLO CLOCKI, E0201, L0701, M0801, DADLO CLOCKI, E0201, L0701, M0801, DADLO	Detail description of logic OCS NA UPD, CAD INPUT SIGNALS NAND GATE LOCATION OF SYMBOL (GS IN PAGE 01) OUTPUT SIGNAL PAGEO 3 4 5 6 F UDP CAD GOSO1 H



T-42-11-09

48

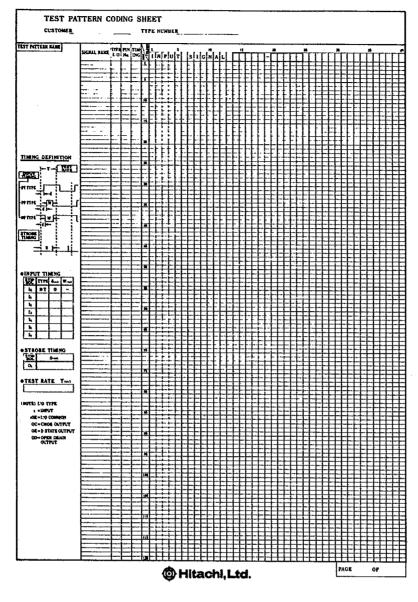
5. GENERATING TEST VECTORS

	Test vector sheet	
5.2	Timing definition	51
5.3	Translation of waveforms	53
5.4	Description of test vectors	55
5.5	Loops	56
5.6	Restrictions on test cycles	57
5.7	Coding example	58

www.DataSheet4U.com 5.1 TEST VECTOR SHEET

T-42-11-09

It is recommended that test vectors are described on Hitachi format coding sheets. Coding sheet:



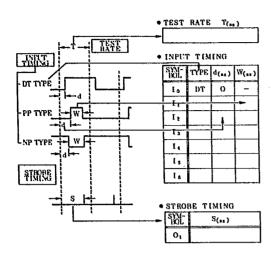
At first a user should define the timing required for testing. The test equipment provides 7 timing generators, I_0 through to I_6 . I_0 is fixed as DT type signal with d= 0. Further, only one strobe timing is available for each set of test vector.

-timing restrictions-

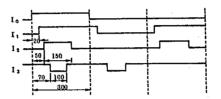
 $T \ge 150$ $T-(d+w) \ge 20$ $d \ge 20$ T-S ≥ 30

W ≥ 50

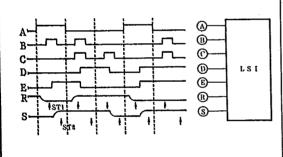
\$ ≧ 20



TEST RATE		300ns			
	SYMBOL	TYPE	d(ns)	w(ns)	
	I ₀	DT	-	-	
	I ₁	DT	20	-	
	I2	PP	50	150	
INPUT	I 3	NP	70	100	
	14				
	I ₅				
	16			<u> </u>	
OUTPUT	01		250ns		



<Example 1>



In the example, B has the same timing as C while D as E. Consequently 3types of timing are required. The timing for output R is different from S and requires different strobe timing.

For this situation, test vectors need to be run twice with S1/S2 strobe timing respectively. S1 should be positioned so that there is a good margin before output S changes.

I/O DATA BUS Negative latch cLK In the actual system, the timing for I₁ is different from I₂ and consequently the timing of O₁ differs from that of O₂. However, during testing they may not necessarily need to follow the actual operation. Such considerations would reduce the complexity of test vectors.

T-42-11-09

Clocks which have no delay are not acceptable, nor is the situation of $d + w \ge T$.

In the actual testing, I/O switching occurs at To. So, latch clock should be shifted back-wards in order to improve the hold time.

-I/O switching problem

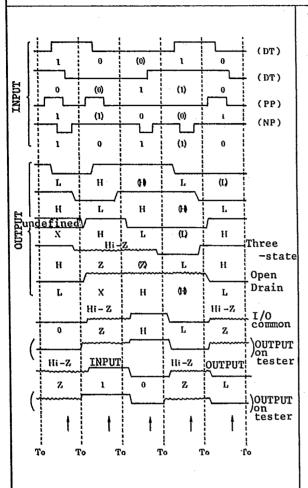
 T_0

Describe the inputs ... 0, 1 outputs ... H, L, Z, X

Input	DT	0, 1
	PP	0, 1
	NP	0, 1
OUTP	UT	H, L, Z, X

Note:

- 1. PP/NP: Pulse is given when "1"
- 2. Z; Hi-Z, X; undefined or masked
- Blank: continue the previous state



* Input pattern; describe the state after a change occurs

DT 0-1 ... 1 1-0 ... 0 no change ... blank

PP/NP Pulse given ... 1 no pulse ... 0 no change ... blank (repetitive)

Output pattern; describe the state at the instant the strobe occurs.

Strobe timing should be set with sufficient margin.

Strobe after the signal transition should be determined after consideration of the maximum delay, on the other hand minimum delays should be considered with reference to the period before the transition.

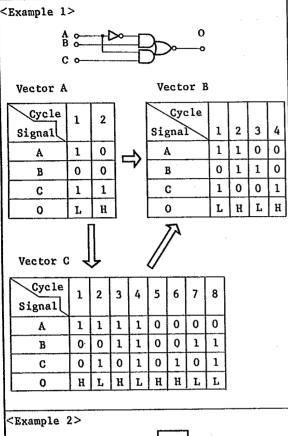
Open drain; when disabled the state should be described ''H'' (transition from an ''L'' to an ''H''), requiring a period of time before stability is achieved depending on the time constant determined by external components. In this situation the state right after the transition should be masked by an ''X''.

* I/O common

input mode ... 0, 1 output mode ... X, H, L, Z

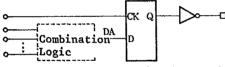
* I/O switching occurs at To in the testing hardware. Thus, if the test vectors are designed with some delay time, both input/output shall be enabled simultaneously.

As mentioned in the previous sections, test vectors are utilized not only for logic verification but also for final testing in mass production. Therefore high fault coverage is required. In order to prepare the sophisticated test vectors, consideration of cell-level fault coverage is not as straight forward as describing vectors to represent the operation of the system.



* If vector group A is just for verifying the function and group C provides all the combinations of inputs; from the standpoint of fault coverage cycle 3 is omittable since it is equivalent to cycle 1.

Thus, vector group B would be the most efficient vector set.



Output of combination logic, DA, needs to be clocked out to an output of the GA.

As mentioned in 3.2 injected fault is ultimately detected at an output of the GA. If many stages of F/F type are involved before the signal appears at the terminal, a very sophisticated set of vectors may be needed.

<Example 3>

Expected outputs should be described together with the associated inputs.

Expected outputs can be compared and checked with the simulation results done by DA program. Expected outputs should be included in test vectors supplied by the customer.

Repetitive patterns can be compressed into a short number of loop cycles where the maximum length of a loop cycle is 4096. During loop cycles, if it is difficult to describe the expected outputs, they can be masked (X).

Expression

Column

LOOP

Loop of limited columns

Loop of block

LOOP 2-4/15 Nesting is not supported.

•Loop of signal block

Vectors

Loop n

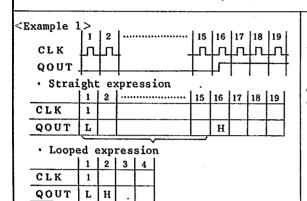
• Loop of limited columns

Vectors

Loop S1-e1/n1, S2-e2/n2

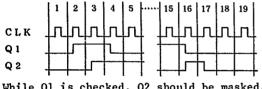
n; loop times, Sn; start column,

en; end column



In this example, 1 through 15 steps both inputs/outputs are repeated. If a looped expression is employed all the cycles described in 15 columns are compressed into 1 column.

<Example 2>



While Q1 is checked, Q2 should be masked.

	1	2	3	4	5	6	7
CLK	1						
Q 1	L.	Н		L			
Q 2	Х				L		

LOOP 1-1/15

LOOP 1-4/4

An output which changes regularly is checkable while the loop is cycled. Outputs that change in an irregular manner should be masked (X). In this case, fault coverage may be degraded.

<Example 3>

Loop \$1-e1/n1, \$2-e2/n2,

Loop Sm-em/nm

When loop commands cannot be described in one line, they can extend to multiple lines with 'Loop' commands at the head of each line respectively.

Test vectors should be designed to be as short as possible in order to achieve a fast throughput on the simulation. Also as test vectors are utilized for production testing, memory capacity imposes a limitation on the number of cycles. In the logic design phase, this limitation should be considered when designing the test circuit.

- * Maximum vector set is 10 and every set requires initialization.
- * The maximum number of cycles for each set including expanded loop cycles is 4000.
- * The total number of cycles before expansion is 4000.
- * Maximum test time is 100 ms.

<1>

The DC parametric test is generated by the DA, which supports up to 2000 cycles. Consequently, all logical states should be covered within 2000 cycles.

e23

While a vector is in progress, changes in input timing definition are not supported. If an input should change, it is necessary to separate vector.

<3>

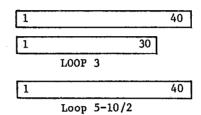
Each individual set of vectors should provide initialization cycles at the beginning.

Timing definition is also required respectively.

DC parametric test: V_{OH} , V_{OL} , input/output leakage current

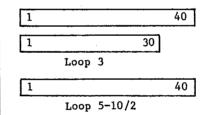
Calculation of cycles required for a vector:

Ex. 1



Cycles: 40+30×3+4+6×2+30=176

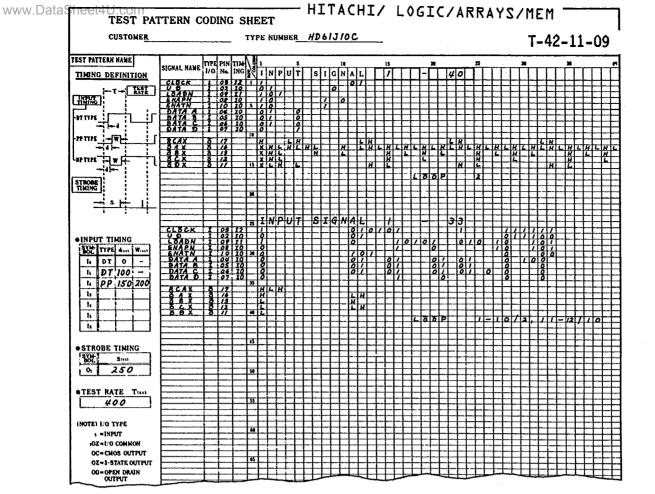
Ex. 2 Total cycles



Total cycles.

40 + 30 + 40 = 110

(Before expanding loop cycles)



6. HITACHI DOCUMENTATION

5.1	Estimation of gate count	60
6.2	Path delay check sheet	61
6.3	I/O transition timing check sheet	62
6.4	Logic drawing sheet	63
	Test pattern coding sheet	
	Machine drawing example	
6.7	Pin assignment example	66
6.8	TTL to Gate Array conversion	68

			Equ.	Shee	t No.	Shee	t No.	Shee	t No.	Sheet	No.
Mac	rocell		count	n pcs	Gate count G × n	n pcs	Gate count G × n	n pes	Gate count G × n	n pcs	Gate count G × n
		1	(G) 1		G × n	Pea	G × n	PCO	G × n	PCS	G × n
Power	NAP	2	2								
gate	NRP	3	3			· · · · · · · · · · · · · · · · · · ·					
	4		4								-
		ι	0.5		-						
		2	1								
		3	1.5							-	
		4	2				-				
Gate	NA	6	9.5								
545	NR	8	5 5		_						
		9	6. 5		1						
		12	8								
		16	1 0. 5								
	EOR/ENOR		2. 5	-							
	LD		3.5								
	LDC1		4								
	LDPC0		6								
Flip-	FD		6								
flop			8								
			6								
			9								
	FJPC1		12								
Latch	LRS0/LRS3		3								
	LR2S20/LR2	S23	4								
S/R	ZSR		10				İ				
	ZSRCP3		12								
	NRA 23/NA R	23	1.5		<u> </u>						
	NR2A2/NA2I	₹2	2		<u> </u>						-
Complex	NR4A2N		4. 5		<u> </u>	ļ. <u></u>					
gate	NR8A2N		9. 5		<u> </u>						
	NR2A3N		3.5								
	NR2A4N		4.5		<u> </u>	<u> </u>	<u> </u>				
Multi-	M2TIN		3						ļ 		
	M4T1N		8.5		ļ		_		<u> </u>		
	M1T2N		3.5	ļ	<u> </u>	<u> </u>	<u> </u>		ļ	<u> </u>	ļ
Decoder	D2T4N		8	ļ	<u> </u>			ļ			
Three-	D3T8		12		ļ				<u> </u>	ļ	
state	ANZ		3. 5	<u> </u>	_		 	ļ	ļ	ļ	<u> </u>
gate	NAZ		1.5	ļ	 	•	ļ	ļ	· .		<u> </u>
	ZEQC4		12	ļ	 	ļ	_	ļ		 	
Others				ļ	 					ŀ	_
			ļ		ļ			ļ			ļ
Total			 			ļ	<u> </u>				
Total	Ļ		<u> </u>	1	<u></u>	L.,	<u> L</u>	L	<u> </u>		<u> </u>

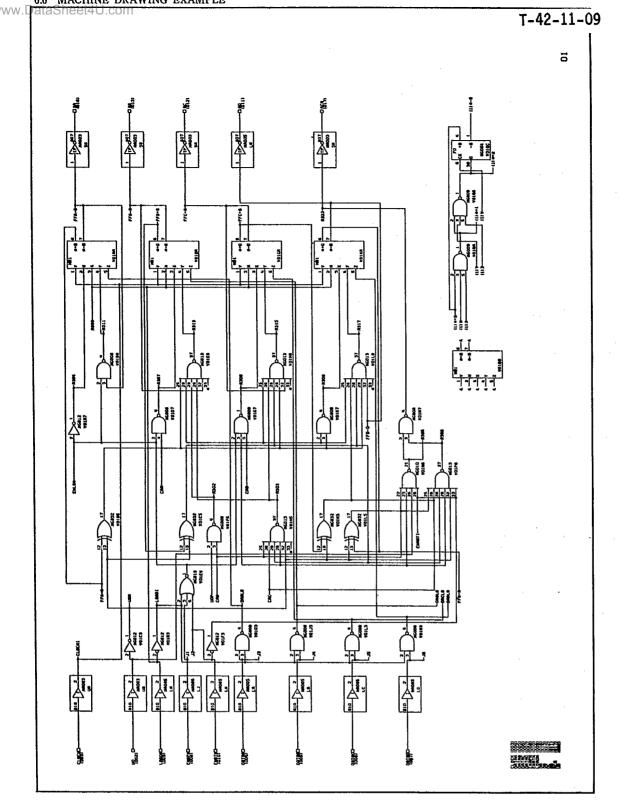
Æ.	CRITICAL PATH/ DELAY CHECK PATH	EQUEST	(typical) ESTIMATE	AFTER ROUTING
1				
2				
3				
4				
5				·
6				
7				
8				
9				
10				
11			•	
12				
13				
14				
15				
16				
17				
18				
19				
20				
Exa- mple	CK _~ D10\		2 1 0	
Note	CK CK +Q CK +Q D Loading capacitance of output pin : 130pF	D10		

IN Æ	PIN Name	I/O Type	тп	MING	PIN	PIN Name	1/0	Type	TIMING
1	- Traine				41	Tranc.			
2		1		,	42				
3					43				
4				• • • • • • • • • • • • • • • • • • • •	44		-		
5					45				
6					46				
7					47				
8					48			- 1	
9					49				
10					50				
11					51	•			
12					52				
13					53				
14				. ·	54				
15	- 				55				
16					56				
17					57				
18					58				
19					59				•
20	,				60				
21					61				
22					62				
23					63				
25					64				
26	 				65				· · · · · · · · · · · · · · · · · · ·
27	· · · · · · · · · · · · · · · · · · ·				67				
28					68				
29				•	69				
30					70				
31	-				71	 ,			
32					72	·····			
33				-	73				
34					74	·			
35					75	-			
36					76				
37				·· · · · · · · · · · · · · · · · · · ·	77				
38					78				
39					79				
40					80				
ſ		Ι0	I 1	I 2	18	I 4		I 5	I 6
	Type	DT			1				
	D (ns)	d = 0							
	W (ns)	-							
	(I) No.		D Inverting =						/0 common = :

62

63

CUSTOMER	TERN C				TYF	E I	NUM	(BE	R																										
			_						Ξ														_				_								
ATTERN NAME	SIGNAL NAME	TYPE 1/O	PIN No.	TIM	\ <u>\$</u>	1	ul p	-lii	5 T		e l 1	Te.	16 N	<u> </u>	F- F	1	1		7-	*		Ŧ	T	×	Ŧ	Т	$\overline{}$	*	Ť	Т	Т	'35 	П	ĊΤ	
					1	1	1	Ľ	H	H	+	ť.		1	出	\pm	\pm	H	-	┢	\exists	\pm	╁	Н		+	\pm		1	\pm		Ł	Н		_
					Н		\pm	\pm	Н	廿		上	H	土	H		\pm		\pm	1		\pm	t	Н		+	H		\pm	士	士	E	Н	Н	
		_	_		3		_	\perp	ᅥ	士	\pm	\pm	Н	土	H	_	±	Н	⇉		╛	士	士	H		+	\pm	Н	\pm	\pm	土	上	Н	ᆸ	
		_				-	1	L		吐	+		Н	\pm	Н	١.	1.	Н	_				\vdash	Н		1	\pm	Ħ	1	\pm	土	£	H	Н	\exists
			\vdash			1	+	Ł	Н		+	\perp	H	士	님	1	1	Н	+	\pm		-	t	Н		_	\pm	Н	\pm	\pm	\pm	土	Н	Н	-
			<u>-</u>		Į\$	\exists	1	1	Н		_	+	Н	\pm		\pm	\pm	Н	\pm	1	Н	-L	Ł	Н	\pm	\pm	\pm	Н	\pm	\pm	\pm	£	\mathbf{E}	Н	\exists
							+	1	Е	<u>-</u>	+		Н	\pm	Н	1	1	Н	+	Ł		\pm	ŀ	Н	\exists	\pm	\pm	Н	\pm	\pm	\pm	Ε	oxdot	\Box	\equiv
					15		1	Ł	Н	\pm	\pm	上	Н	士	Н	\pm	上	H	\pm	士	Н	\pm	Ł	Н		\pm	上	Н		土	\pm	Ε	Đ	Н	
			<u> </u>				<u>- -</u>	Ł	Ы	H.	+	1	Ŀŧ	-	出		±	Ŀŀ	_±	+-	H	1	Ŀ	Н			\pm	Н	\pm	\pm	\pm	E	\blacksquare	Н	_
							Ŧ	1.		Н	1	上	Н	\pm	Н	1	1	Н	1	\pm	Н	\pm	£	L	\exists	1	£	Н		\pm	\pm	\pm	E	Н	_
			-		*	\dashv	Ŧ	F	H	Н	\pm	F	H	Ŧ	Н	\pm	E		-	Ŧ	Н	Ŧ	F	F	Н	\mp	F	H	\pm	Ŧ	Ŧ	F	P	Н	Ξ
No Dominanou		Ш				-1	1	-		Н		1	出	\pm	Н		\pm			1	Н		E			E			± 1	\pm	\pm	E	E	Н	Ξ
NG DEFINITION		1	\vdash		B		1	H	Н	Ш	\pm	土	Н	土	oxdot			Н		Ŧ	Н	\pm	Ε	E		-E	\pm		\pm	\pm	\pm	£	\pm	Н	=
T [[]		E			Е	\mathbf{J}	\mathbf{F}	£	E	H	E	E	H	\pm	Ø	\pm	£	Ø	\pm	F	Ø	Ŧ	F	E	Ø	Ŧ	F	П	J	Ŧ	Ŧ	£	£	Ø	Ē
5		Е			E	\exists	F	E	E	H	£	壬	B	壬	Ð	\exists	Ŧ	П	\pm	Ε	Ø	Ξ	F	E	Ы	\pm	£	Ø	J	\pm	Ŧ	£	£	В	Ē
<u>"- -</u> -		Е	E	E	39	Ы	Ŧ	E	E	H	\pm	E	\Box	\pm	\Box	\pm	F	E	\pm	E	\exists	\pm	F	E	В	\pm	£	В	\exists	\pm	Ŧ	£	F	Н	Ē
##T		E		E	E	oxdot	\pm	E	E	H	Ξ	£	H	\pm	Ð	H	F	Н	\pm	F	Н	\pm	Ŧ	E	Ы	\pm	£	Н	J	\pm	£	£	£	Н	Ē
r - V		E	E	E	×	H	\pm	F	E	H	F	F	₽Ŧ	F	F	Ŧ	Ŧ	Ø	\pm	£	Ø	\pm	E	E	В	Ŧ	F	F	H	Ŧ	Ŧ	£	£	Ð	Ē
	 		E.,	F	Е	Н	\top	F	E	Н	Ŧ	F	H	Ŧ	H		Ŧ	П	\blacksquare	Ε	Н	T	F	F	Н		∓	H	H	7	Ŧ	Ŧ	F	P	F
		\vdash	-			H	Ŧ	+	-	Н	7	F	H	Ŧ	\exists	F	Ŧ	F		Ŧ	H		F	F	Н	7	Ŧ	Н	Ħ	Ŧ	Ŧ	Ŧ	Ŧ	П	F
			F	E	4	Н	-	F	F	П	Ŧ	Ŧ	H	Ŧ	Π		Τ	Н		Ŧ	Н	\mp	F	F	Н	7	+	H	H	干	Ŧ	F	F	F	F
			-	F		H	Ŧ	F	F	П	-	F	H	干	Ħ		F	F	-	F	FI	-	Ŧ	F	H	7	Ŧ	H	H	7	Ŧ	Ŧ	Ŧ	F	F
, , ,			F	-	45	Н	-	F	F	Н	-	F	П	干	\sqcap	Ŧ	Ŧ	F	-	+	П	7	F	F	П	7	Ŧ	П	Н	7	Ŧ	Ŧ	F	F	F
 		Е	F	-	П	П	Ŧ	F	F	П	7	Ŧ	П	Ŧ	\sqcap		Ŧ	П	1	Ŧ	П		Ŧ	F	H	7	+	F	Ħ	#	Ŧ	Ŧ	F	F	F
					F	П	\mp	F	F	П	Ŧ	Ŧ	H	Ŧ	Ŧ	\blacksquare	Ŧ	F	\mp	Ŧ	FI	-	F	F	H	7	Ŧ	F	П	7	Ŧ	+	Ŧ	F	F
			-	F	50	H	-	F	F	Н	7	Ŧ	П	\mp	Ħ	\exists	+	F	\exists	Ŧ	Н	7	Ŧ	F	П	7	Ŧ	F	H	7	Ŧ	Ŧ	Ŧ	F	F
TYPE 411 Was			F			П	7	Ŧ	F	Н	Ŧ	Ŧ	Ħ	7	H	H	-	F		Ŧ	П		F	F	П	7	Ŧ	F	H	7	Ŧ	Ŧ	F	F	F
DT Q -		=	\vdash	F	55	H	7	Ŧ	F	П		Ŧ	Ħ	-	Ħ	H	Ŧ	F	7	F	П	7	Ŧ	F	П	Ŧ	Ŧ	H	Ħ	7	丰	Ŧ	1	F	F
101 0 -		F	F	-	-	П	Ŧ	Ŧ	F	П	Ŧ	Ŧ	П	Ŧ	H	F	Ŧ	F		+	H	7	Ŧ	F	П	7	Ŧ	H	Ħ	7	Ŧ	Ŧ	Ŧ	F	F
1-1-1-1				F	Е	П	Ŧ	F	F	П	Ŧ	Ŧ	П	\mp	F	\Box	Ŧ	F	\Box	Ŧ	П	-	F	F	H	+	Ŧ	H	H	7	Ŧ	Ŧ	F	F	F
			F	\vdash	60	Н	Ŧ	Ŧ	F	Н	Ŧ	Ŧ	H	\mp	F	H	Ŧ	F	\sqcap	\mp	F	-	Ŧ	F	П	7	Ŧ	F	Ħ	7	Ŧ	Ŧ	Ŧ	F	F
		\vdash	\vdash	\vdash	F	П	-	i	F	\Box	\mp	Ŧ	Π	Ŧ	\Box	Н	-	F	H	-	H	-	Ŧ	F	П	1	丰	F	П	7	丰	Ŧ	Ŧ	F	F
			\vdash	\vdash	65	Н	\pm	\pm	F	Н	-	Τ	Н	-	P	H	+	F	H	-	Н	-	Ŧ	F		\equiv	\mp	F	П	\mp	Τ	Ŧ	F	F	F
			E	Е	E	Н	\pm	+	F	Н	\pm	\pm		\pm	\pm	Н	\pm		\Box	-E	Е		\pm	Ε		\pm	\pm	Е	\Box	\pm	Ξ	Ξ	Ε	E	F
					E	Н	1	1	Ł	Н	\pm	土	H	\pm	H		\pm	E	Н	-	Н	-	+	\vdash		\pm	土	E	H	\pm	Ŧ	Ŧ	E	Е	E
ROBE TIMING		\vdash	L	\vdash	10	Н	_	\perp	L	Н	1	\pm	Н	\pm	\pm	Н	L	E		L	Н	_	\pm	£	Н	-	\pm	E	\Box	\pm	\pm	\pm	\pm	E	E
Sun		L		上	┢╌	Ы	\pm	+	Ł	Н	\pm	土	Ы	土	上	\vdash	_	Ŀ	Н	\perp	Н	\pm	\pm	£	Н	-	\pm		$oxdapsymbol{oxdapsymbol{eta}}$		\pm	Ξ	\pm	E	F
		Е		Е	15	Н		E	Ŀ	Н	-	Ŧ	\Box	\pm	F	Н	Ŧ	F	Н	Ŧ	Η	$\overline{}$	-	F		\blacksquare	Ŧ	Н	П	\mp	Ŧ	Ŧ	Ŧ	F	F
PT DATE T		-	-		\vdash	-	-	+	H	Н	+	Ŧ	H	\pm	\vdash	Н	\pm	F	H	Ŧ	Η	T	\pm	E	Н		Ŧ	F	H	\exists	Ŧ	Ŧ	Ŧ	F	F
ST RATE Ton				E	L	Н	1	Ĺ	È	Н		\pm	\Box	\mathbf{H}	\perp			E		\pm	H	\pm	\pm	\pm	Н		\pm	Ł	Н	\pm	Ξ	Ξ	Ξ	Е	E
	-	\vdash	-	+	**	Н	-	+	╁	Н	+	+	Н	\pm	+	Н	+	\vdash	Н	+	-	+	Ŧ	+	Н	-	\pm	F	Н	\mp	Ŧ	Ŧ	Ŧ	F	F
E) I/O TYPE		E	\vdash	E	E	П	Ŧ	I	£	H	Ŧ	Ŧ	F	H	£	\Box	£	E	H	F	E	F	F	F	E	H	Ŧ	F	Н	J	Ŧ	Ŧ	F	F	F
i =INPUT		E	oxdapprox		15	\Box	\pm	F	F	\Box	Ŧ	Ŧ	F	EF	F	H	\pm	F	H	Ŧ	E	\exists	Ŧ	F	E	F	Ŧ	E	H	7	Ŧ	F	F	F	F
OZ=T/O CONNON		E	E	E	E	Ø	Ŧ	Ŧ	F	Ø	Ŧ	Ŧ	Ð	F	\pm	Ø	F	F	Ħ	F	F	F	Ŧ	F	F	F	Ŧ	F	П	7	Ŧ	Ŧ	Ŧ	F	F
OC=CHOS OUTPUT OZ=3-STATE OUTPUT		E	E	E	E	Ø	\pm	Ŧ	F	Ð	Ŧ	Ŧ	\Box	F	F	П	F	F	H	1	F	-	Ŧ	F	E	H	Ŧ	F	П	7	Ŧ	Ŧ	F	F	F
OD=OPEN DRAIN		E	F	F	190	Ø	\exists	Ŧ	F	F	J	Ŧ	F	H	F	Ħ	Ŧ	F	H	Ŧ	F	7	Ŧ	F	F	H	Ŧ	F	F	7	+	Ŧ	Ŧ	F	F
OUTPUT		E	F	F	F	Ħ	1	Ŧ	F	F	-	Ŧ	F	H	F	П	Ŧ	F	H	T	F	\exists	Ŧ	F	F	H	Ŧ	F	Ħ	7	7	Ŧ	Ŧ	F	F
•		E	F	1	95	F	_	F	F	Н	Ŧ	Ŧ	\mathbf{F}	F	\pm	H	Ŧ	F	H	Ŧ	F	H	Ŧ	T	F	H	Ŧ	F	F	7	7	丰	丰	F	F
			\vdash	\vdash	F	Н	1	Ŧ	F	П	\exists	Ŧ	F	F	Ŧ	H	Ŧ	F	H	-	F	H	Ŧ	F	F	Ħ	Ŧ	F	F	\dashv	Ŧ	7	Ŧ	F	É
		F	+	F	F	H	\dashv	+	Ŧ	П	\dashv	Ŧ	=	FF	Ŧ	H	+	F	H	- -	-	-	+	1	Е	H	#	F	Ħ	Ħ	#	‡	+	F	F
		F	\vdash	F	100	F	7	7	Ŧ	\Box	-	+	P	H	丰	H	7	F	П	+	F	\dashv	+	7	F	\dashv	+	-	Ħ	H	#	+	+	Ŧ	F
		F	+	-	F			+	F	Н	-	+	F	H	Ŧ,	H	-	-	П	+	F	Ħ	7	Ŧ	F	Ħ	丰	F	Ħ	\dashv	7	‡	+	Ŧ	F
		T	\top	F	105	H	Ħ	#	1	Ħ	H	#	Ŧ	FF	Ŧ	H	+	F	H	+	F	Ħ	#	+	F	Ħ	#	F	Ħ	Ħ	#	#	‡	Ŧ	F
		F	F	\vdash	-	F	Ħ	+	T	Ħ	7	#	F	F	+	Ħ	1	1	Ħ	‡	-	H	+	+	F	ΓÍ	#	F	Ħ	Ħ	#	‡	‡	#	F
		F	1	1	F	H	Ħ	#	+	Ħ	Ħ	#	#	岸	#	Ħ	#	#	Ħ	#	F	Ħ	#	+	Ħ	H	#	‡-	Ħ	Ħ	#	#	#	#	F
		#	+	丰	110	F	Ħ	#	#	H	\vdash	#	#	炐	#	Ħ	#	#	Ħ	#	t	Ħ	#	#	+	H	#	+	Ħ	H	#	#	#	#	t
		1	1	1	+	F	Ħ	+	+	Ħ	\exists	#	+	Ħ	#	Ħ	+	#	Ħ	_	1	Ħ	+	‡	+	Ħ	士	士	Ħ	H	#	#	#	#	t
		1	+	+	1.		H	#	‡	Ħ	Ħ	#	+	井	#	H	+	+	Ħ	\pm	1	Ħ	#	#	F	H	#	t	Ħ		#	#	#	‡	t
		\vdash	1	1	111		Ħ	\pm	#	Ħ	Ħ	#	+	片	+	H	#	+	Ħ	\pm	\pm	Ħ	\pm	+	t	H	#	#	Ħ	\exists	\pm	#	#	#	t
		\pm		\pm	+	H	Ħ	\pm	+	±	H	#	\pm	H	\pm	Н	#	+	Н	\pm	t	Н	+	±	L	Н	\pm	\pm	${}_{\!$	Н	\pm	\pm	\pm	\pm	t
	J	+-		4	130	\mathbf{L}	 -	_	_	ட	ட	_1			1.	1 . I	- 1	1	ı [1	1	ı. I	_1_	Ľ	1	ιſ	- 1 -	L.	1.]	ı [[ட	1	Ĺ



Type No.

Customer:

T-	42	-11	-09
----	----	-----	-----

PIN No.	PIN NAME	I/O type
1	GND	
2		
3		
4		
5		
6		
7		
8		
9		
10		
11		
12		
13		
14	v _{cc}	

PIN No.	PIN NAME	I/O type
15	GND	-
16		
17		
18		
19		
20		
21		
22		
23		
24		
25		
26		
27		
28	v _{CC}	_

Note: Power supply pins are fixed.

I/O types

i : Non inverting input buffer

iOZ : I/O common

OC : CMOS inverting output buffer

OZ : Three-state output buffer

OD : Open drain output buffer

4 M 2

HD61J×××P

JAPAN

Product version: HD61K

Customer

Package type FP80 pin

Type No.

PIN No.	PIN NAME	I/O type
1	ИС	
2		
3	NC	_
4		
5		
6		
7		
8		
9		
10		
11		
12		
13		
14		
15		
16		
17		
18		
19		
20		
21		
22	NC	
23		
24	NC	
25		
26		
27		
28		
29		
30		
31	v _{CC}	
32		
33	GND	
34		
35		
36		
37		
38		

PIN NO.	PIN NAME	I/O type
41	NC	T - T
42		
43	NC	_
44		
45		
46		
47		
48		
49		
50		
51		
52		
53		
5 4		•
55		
56		
57		
58		
5 9	· · · · · · · · · · · · · · · · · · ·	
60		·
61	·	
62	NC	
63		
64	NC	
65		
66		
67		
68		
69		
70		
71		
72	v _{CC}	
73		
74	GND	
75		
76		
77		
78		
79	1	-

Notes: Power supply pins are fixed.

I/O types

39

40

i : Non inverting input buffer

iOZ: I/O common
OC: CMOS inverting output buffer
OZ: Three-state output buffer
OD: Open drain output buffer
NC Pins should not be utilized.

◆ 4M2 HD61K×××F JAPAN

80

Notification on conversion

The following conversion table indicates the equivalent gate count when one of the built-in functions is utilized.

For instance HD74LS139 integrates dual 2-line to 4-line decoders, then equivalent gate count 8.5 represents one of the two decoders.

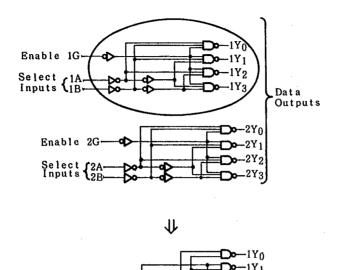
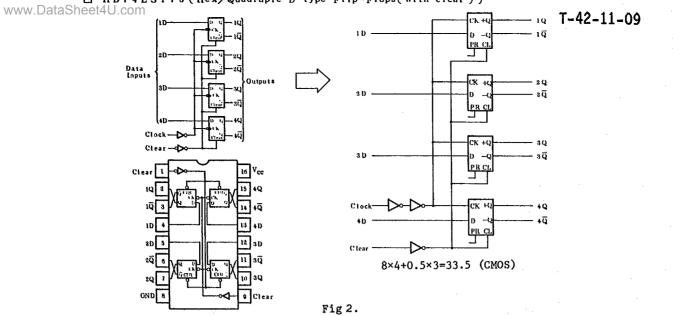


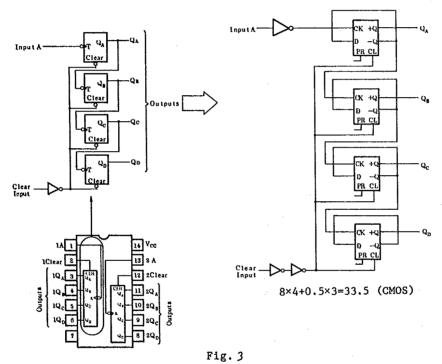
Fig. 1

Equivalent gate=5

This example shows that we have the possibility to reduce the gate count in reality if we do not utilize all pin functions.



☐ HD74LS393 (Dual 4-bit Binary Counters)



70

SECTION II: HD61 Series Product Specification

71

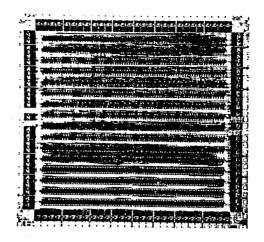
HD61J/HD61K/HD61L/HD61MM

The HD61 series is a master slice CMOS gate array using 2-layer metal interconnect technology. This series has four master chips with gate counts ranging from 504 to 2496, and I/O terminal counts ranging from 50 to 104. These chips can replace not only CMOS logic but also TTL logic due to their high speed of 3.5 ns typ and the compatibility of input and output buffers at TTL level.

LSI design is fully automated by DA (Design Automation) system and a custom LSI is developed based on logic diagram, timing and electrical specifications and test patterns from customer.

FEATURES

- Technology
 Silicon-gate 3-micron CMOS technology
 Dual-layer metal Interconnection
- Low power dissipation
 At 10MH_z operation (Internal gate) 130μW/gate typ
- Abundant input and output configuration
 Allocation of all pins except power supply pins to input/output/input-output
 Output can be CMOS/open drain/3-state
- Memory on-chip (HD61MM)
 Flexibility of memory capacity and word organization
 Selection of single port / dual port memory



- Wide operating temperature range -20 to +75°C
- Wide package selection
 Featuring plastic packages with high pin counts
- Powerful design support
 User-Defined-Macro
 Worldwide network of design and engineering support centers
 - Logic simulation
 - Fault grading
 - Timing verification (delay simulation)

LINE UP

			HD61J	HD61K	HD61L	HD61MM
Gate cour	ıt		504	1080	1584	2496
I/O pin co	ount		50	68	68	104
RAM on	chip		-		-	16 bits x 32 words max
		DP28	0	0	0	<u> </u>
	DP40 DP42	DP40	0	0	0	-
	မှု 🗌	DP42	0	0	0	
	Plastic	DP64		0	0	0
Dankana	~ [FP54	0		_	
Package		FP80		0	0	_
		FP100		_	_	**
	0	DC28	0	0	0	**
	🖁 🗌	DC40	0	0	0	**
	DC40 PGA72	PGA72	-	_	0	
	$I \check{} \Gamma$	PGA120		_	_	0

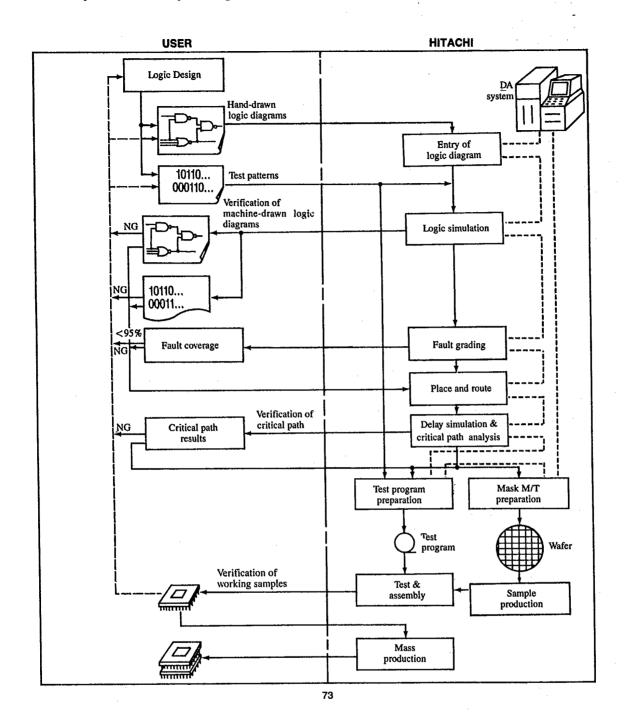
^{**}Under Development

The gate array design flow is shown below.

The interface between the user and Hitachi is based on hand-drawn logic diagrams and test patterns provided by the user.

Hitachi will input these into our computer for logic simulation,

fault grading, placement and routing, delay simulation and test program preparation. After verification of above by user, samples will be produced, checked by user and mass production will begin.



■ ABSOLUTE MAXIMUM RATINGS

	Item	Symbol	Rating	Unit
Supply Voltage*		V _{CC}	-0.3 to + 6.7	V
Terminal Voltage*			-0.3 to $V_{CC} + 0.3$	V
Operating Temperature		Toor	-20 to + 75	°C
Storage	With Bias	Tbias	-20 to + 85	°C
Temperature	Without Bias	T _{stg}	-55 to +125	°C
Output Current Per Output Pin			-8 to 8	
	Total	IOH	-40 to 40	MA

*With respect to GND

Note) Permanent damage may occur if maximum ratings are exceeded.

Normal operation should be under recommended operating

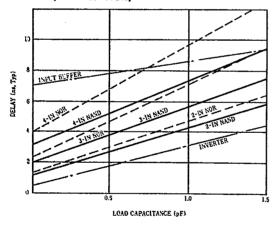
condition, that is GND \leq (Vin and/or Vout) \leq V_{CC}. If these conditions are exceeded, reliability of LSI could be affected.

■ ELECTRICAL CHARACTERISTICS (Vcc=5V ± 5%, Ta=-20 to +75°C)

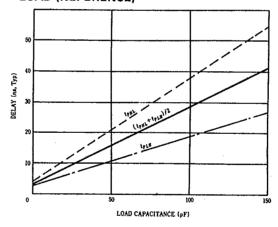
	Item	Symbol	Test Conditions	min	typ	max	Unit
Input Voltage		VIH		2.2		Vcc+0.3	V .
Input Vollage		v_{IL}		-0.3	_	0.8	V
Output Voltag		V _{OH}	$I_{OH} = -2mA$	2.4		-	V
		V_{OL}	$I_{OL} = 2mA$			0.4	V
Input Leakage Current		I _{LI}	VIN = GND or VCC		i -	1	μΑ
Output Leaka	ge Current	I _{LO}	at high impedance		_	1	μA
	Internal		2 input NAND, FO = 3, Metal = 3mm		3.5		ns
Gate Delay	Input Buffer	t _{pd}	FO = 3,Metal = 3mm	1 -	9	1 -	ПS
	Output Buffer	•	$C_L = 50PF$		24	1 =	ns
Power Dissipa	sipation I _{CC} Internal 2 input NAND, 10MHz - 130		-	μW/Gate			
Terminal Cap	Terminal Capacitance*		$Ta = 25$ °C, $f = 1$ MHz $V_{iN} = 0$ V	—		12.5	pF

^{*}This parameter is sampled, not 100% tested.

■ INTERNAL GATE DELAY vs LOAD (REFERENCE)



■ OUTPUT BUFFER DELAY vs. LOAD (REFERENCE)



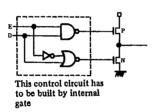
1. I/O BUFFERS

www.DataSheet4U.com

7	facrocell facrocell			Clamp				Delay		
Function			Normalized	Level	Symbol	Symbol	tPLH	(ns)	tPHL	(ns)
Macro Function Name	Equivalent Circuit	lent Gate Count	Load Factor	when Open	3 y 111001	No.	tolh	kLH	tOHL	kHL
Input Buffer	6x0	-				DI	7.7	2.0	6.7	1.0
I/O Buffer	*** <u>-</u> - <u>-</u> - <u>-</u> - <u>-</u> - <u>-</u>	· _	2.2 3.2		P P	D2		put Se nput I	ee Buffer'	,
ANIO	620 - 5h 5h 4g				+1			utput : -state	See Buffer	,,
3-state Buffer	CXD ————————————————————————————————————	-	2.2 3.2		<u>Ç</u> P	D1	8.0	0.12	22.0	0.19
Output	Vcc — CND — CND	_	4.7			D1	8.5	0.12	23.5	0.19
Open Drain Output	CXB		3.2			D1	_		22,0	0.19
NAOD			1.							

Note) Application of Tri-state buffer





2. POWER GATES

Ma	crocell	, i		<i>(</i> 1)	· · · · · · · · · · · · · · · · · · ·			Delay		
Function		Equiva- lent Gate	Normalized Load	Clamp Level	O11			(ns)	tPHI	, (ns)
Macro Function Name	Equivalent Circuit	Count,	Factor	when Open	Symbol	Symbol No.		kLH	tOHL	kHL
Power inverter NAP	>	1	1.4	@	-		1.2	1.2	0.9	1,5
Power-2 input NAND NAP	=	2	1.4		<u>_</u>	_	1.3	1.9	1.2	1.3
Power-3 input NAND NAP	≡ D	3	1.4			_	2.1	2.0	1.9	1.7
Power-4 input NAND NAP		4	1.4		-	_	2.7	2.2	2.9	2.1
Power-2 input NOR NRP	—	2	1.4				1.9	2,5	0.9	1.5
Power-3 input NOR NRP	≡>	3	1.4		⇒ ~		•3.7	3.3	1.5	1.5
Power-4 input NOR NRP		4	1.4		P>-		6.1	4.3	1.8	1.5

3. GATES

	crocell	7	Normalized	Clamp	-]	Delay		
Function Macro Function Name	Equivalent Circuit	Equiva- lent Gate Count	Load Factor	Level when Open	Symbol	Symbol No.	tPLH tOLH	kLH		kHL
Inverter NA	→>	0.5	1	@	→	_	0.6	3.0	0.4	2.2
2-Input NAND NA	₽	1	1			_	1,3	3,6	1.0	2.7
3-Input NAND NA		1,5	1		₽	-	2,0	3.9	1.9	3.4
4-Input NAND NA	= D	2	1		₽	_	2.9	4.3	3.2	4.1
6-Input NAND NA	—	4.5	1		-	_	4.1	2.7	6.0	2.0
8-Input NAND NA		5.5	1				4.7	2.9	7.7	2.0
9-Input NAND NA		6.5	1				4.6	2.9	8.2	2.2
12-Input NAND NA		8	1				5.0	3.0	9.9	2.2

	acrocell		ļ	Clamp		Γ	I	Delay		
Function Macro Function Name	Equivalent Circuit	Equiva- lent Gate Count	Normalized Load Factor	Level when Open	Symbol	Symbol No.		(ns)	tPH)	(ns)
16-Input NAND		10.5	1	Open		_	5.3	3.0	12.5	2.5
2-Input NOR NR	=>-	1	1			-	1.8	4.7	1.0	2.7
3-Input NOR NR	≡ D⊶	1.5	1		⇒	- .	3.7	6.8	1.4	2.7
4-Input NOR NR	\$ -	2	1		⇒	_	6.5	8.5	1,5	3.1
6-Input NOR NR		4.5	1			_	7.6	2.6	3,6	2.1
8-Input NOR NR		5.5	1			_	10.5	2.8	3.7	2.0
9-Input NOR NR		6.5	1			_	8.6	2.8	4.2	2.1

Ma	crocell		1	Clamp		[1	Delay		
Function		Equiva- lent Gate	Normalized Load	Level	Symbol	Symbol	tPLH	(ns)	tPHL	(ns)
Macro Function Name	Equivalent Circuit	Count	Factor	when Open		No.		kLH	tOHL	kHL
12-Input NOR NR		8	` 1				12.1	2.8	4.5	2.0
16-Input NOR NR		10.5	1			_	13.2	3.0	4.9	2.1
2-Input EOR EOR	⇒D—	2.5	1.4		#>-	_	4,2	4.5	3.9	2.0
2-Input ENOR ENR	⇒>>	2.5	1.4		#>~	_	3.2	2.5	3,2	3.0

4. 3-STATE GATES

	Macrocell								De	lay			
Function	Paulus	lent	Normal- ized	zed Level	Level Cumbai	Symbol	Symbol			t _{PLH} (ns)		t _{PH1}	, (ns)
Function Name	Equivalent Circuit	Gate Count	Load Pactor	When Open	Symoor	No.	Input Name	Output Name		K _{LH}	ЮНІ	K _{HI} ,	
3-State Inverter (Internal)		1.5	1	#@@		_	D		2.4	2.6	1.8	1.9	
NAZ	Ε					E/Ē		1.0		0,4			
3-State Buffer	Ē	3.5	1.4	#		_	D		2.4	2.3	3.0	1,7	
(Internal) ANZ	· .		1	@	- >-		E		3,9	2.5	2.9		

5. AND-NOR, OR-NAND GATES

	Macrocell			, ,			-		De	lay		
Function	Equivalent	Equiva- lent Gate	Normal- ized Load	Level	Symbol	Symbol			₽ <i>LH</i>	(ns)	t _{PH}	L (ns)
Function Name	Equivalent Circuit	Count	Factor	When Open		No.	Input Name	Output Name	t _{OLH}	K _{LH}	†OHL	K _{HL}
2-AND- NOR	=D-D-	1.5	1			Al	AND Input		2.6	5.5	1.4	2.9
NRA23							NOR Input		2.0	_	1,4	-
2 Wide- 2 Input AND-NOR NR2A2	=D	2	1			A1			3.0	4.1	2.2	2,8
2-OR- NAND	<u>=D-</u> D-	1.5	1	-		Al	OR Input		2.5	5.0	1.9	3.8
NAR23							NAND Input		1.9	_	1.2	-
2 Wide- 2 Input OR-NAND NA2R2		2	1			AI			3.8	5.0	1.8	2.7

	Macrocell							-	De	lay		
Function Function Name	Equivalent Circuit	Equiva- lent Gate Count	Normal- ized Load Factor	Clamp Level When Open	Symbol	Symbol No.	Input Name	Output Name	-	(ns)	[†] PHL	
4 Wide- 2 Input		4.5	1	#######		A4		-Y	8.5	5.1	3.8	2.6
NR4A2N		7.3		####				+Y	5.6	3.5	9.0	3.5
8 Wide- 2 Input AND-NOR		9.5	1	#################		A5		-ч	6.3	2.5	10,5	3.2
NR8A2N) A A A			#######				+Y	11.8	3.4	6.8	2.4
2 Wide- 3 Input	≡ D ₇ , }>-	3,5	1	-		A2		-Ү	4.4	2.9	4.5	3.1
AND-NOR	==D ₁							+Y	6.2	4.0	5.0	2.8
2 Wide- 4 Input		4.5	1			A4		-Ч	4.5	2.6	7.2	3.8
AND-NOR NR2A4N		7.5				1.4		+Y	9.4	4.9	5.1	2.9

6. LATCHES

	Macrocell								D	elay		
Function Function	Equivalent Circuit	Equiva- lent Gate Count	Normal- ized Load Factor	Clamp Level When Open	Symbol	Symbol No.	Input Name	Output Name		(ns)	t _{PH I}	(ns)
Name	Cucait	Count	Pactor	Opt.ii			Pante	1,12,110	[†] ОLН	K _{LH}	t _{OHI}	K _{HL}
RS-Latch		3	1			A3	₹ R	+Q	5.1 3.0	2.5	3.3	1.5
LRS0	2—D—J >>,						₹ R	-Q	3.4 5.6	2.3	2.9	1.5
RS-Latch	5	3	1		<u></u>	А3	S R	+Q	2.9	2.3	4.0 5.5	1.8
LRS3	" → D → B		•			AS	S	-Q	3.0	2.3	5.8 4.2	1.8
R2S2-							\overline{S}	+Q	7.2 4.5	2.6	- 3.4	1.6
Latch	5-50-0-8	4	1			A4	Ī		4.4		3.7	
LR2S20							R	Q	6.9	2.6		1.7
R2S2-	\$ \						S	+Q	3.3	2.3	6.4	2.2
LATCH		4	1	-		A4	R				7.5	ļ
LR2\$23	10.						S R	–Q	3.3	2.3	7.5 6.6	2.2
				@	LD		G	+Q	2.7	2.2	3.1	1.7
D-Latch	- G 4	3.5	1.4	@	b -v	С	D		3,4		3.9	1.7
LD	U Ü						G D	-Q	5.0	2.2	3.6 4.4	1,3
					Luci		G		3.5		3.2	<u> </u>
D-Latch with CLR			1.4	@			CL D	+Q	2.8 4.2	4.3	3.9	1.8
	U	4	1	@		С	G		4.6		4.7	
LDC1			1	#			CL D	-Q	3.3 5.2	2.3	3.6 5.4	1.6
D-Latch with PRE/CLR	eR v	6	1,4	#		С	G PR CL D	+Q	8.6 6.6 4.5 9.1	2.7	5.3 - 3.4 6.8	1.9
LDPC0	D CX		1 1	@@	PR CL		G PR CL D	-Q	9.1 4.5 7.3 10.5	2.7	5.3 3.1 - 5.6	1.8

7. FLIP-FLOPS

	Macrocell	Equiva-	Normal-	Clamp					De			7.3
Function Function Name	Equivalent Circuit	lent Gate Count	ized Load Pactor	Level When Open	Symbol	Symbol No.	Input Name	Output Name	t _{OLH}	(ns) K _{LH}		(ns) K _{HL}
DFF		6	1	0 0	CK +Q	С	CK	+Q	4.6	2.4	3.9	1,6
FD	0 0							-Q	5.2	2.2	5.5	1.4
			1.4	@			СК	+Q	2.5	2,5	3.1	2.1
DFF with Load		6	1.4	@#@@#@	CKN +Q DC DL	B4	DL		7.7		9.0	
FDL			1 1	# @	LN -Q		CK DL	-Q	4.3	2.2	3.6 8.5	1.4
					FDPC3		CK		5,4		4.2	1.7
DFF with PRE/CLR	CK Q	8	1	@	D -Q	c	CL PR	+Q	2,5 5. 5	4.4	1.2	1.7
			1,4	#	PR CL		CK		6.8 5.8		6.8	1.5
FDPC3			1.4	#			PR	-Q	4.0	4.4	1.5	1,5
jkff	J Q	9	1.4 1 1	@@#		C	CK	+Q	3,8	2,3	6.8	2.8
FJ	dx 0		1					-Q	5.0	2.5	5.4	2,3
					FJPC1		CK	ļ	7.2	1	7.5	\
JKFF with PRE/CLR			1.4	@ #	1 CK +0		PR CL	+Q	3.0	2.3	7.2	5.8
FKE/CLK	K CIT OF	12	1	1	PR CL	C	СК		4.7	<u> </u>	10.1	
FJPC1			1.4	@ #			PR	┥ `	4,4	2.3	5.7 4.5	4

8. MULTIPLEXERS

	Macrocell								D	elay		
Function		Equiva- lent	Normal- ized	Clamp Level	Symbol	Symbol			t _{PLH}	(ns)	₹PHI.	, (ns)
Function Name	Equivalent Circuit	Gate Count	Load Factor	When Open	Symbol	No.	Input Name	Output Name	[†] ОLН	K _{LH}	t _{OHL}	K _{HI} ,
							Yo		3.3		5.0	
2 to 1 Multi-	"	3	1.4 1	####	M2T1X S - Y T ₂ - Y	, na	Y ₁	+Y	4,4 3,8	3.2	3.8 4.7	2.8
plexer	" - -D _		1	#		B2	Y ₀		4,4		2.0	-
M2T1N							Y ₁	_Y	3.3	4.2	3.1	2,6
			<u> </u>		· <u>-</u>		_					
4 to 1	" — • • • • • • • • • • • • • • • • • •	•		####	— A A — B — Ye → Y		Y ₀ Y ₁ Y ₂ Y ₃ A	+Y	9.5 8.6 6.2 4.8 11.3	4.8	9.3 9.6 6.8 5.8 15.5	3.7
Multi- plexer M4T1N	" —— - 1 " —— - 1	8,5	1	######		В4	B Y ₀ Y ₁ Y ₂ Y ₃ A	Y	10.5 8.5 8.8 6.0 5.0	6.7	13.0 7.5 6.3 4.6 3.3 8.5	3,5
-							Y	+0	3.4	3.4	7.2 2.6	2.2
1 to 2 Demulti- plexer	Y ————————————————————————————————————	3.5	1.4	#	MIT2X A -1 -1 -1 -1	В3	A Y A Y A	+1	2.7 2.9 2.0 2.8	3.5	3.5 2.3 3.1 2.0 3.1	2,3
MIT2N	<u></u>						Y	-1	1.6	2.3	1.6	2,3

	Macrocell									lay		
Function	Fautralent	Equiva- lent	Normal- ized	Clamp Level When	Symbol	Symbol No.	Input	Output	tPLF	(ns)	t _{PH1}	(ns)
Function Name	Equivalent Circuit	Gate Count	Load Factor	Open		NO.	Name	Name	^t OLH	K _{LH}	^t OHL	K _{HL}
	,,,,,,						A B	-0	3.3	2.3	4.0 3.9	2.3
					·		A B	-1	4.9	2.3	4.7 3.8	2.3
					D2T4N +0		A	2	3.7	2.4	4.1	2,3
2-bit				#	A +1 -1 -1		A	-3	5.1	2.2	4.7	2.3
Decoder		8	1	#	+2 	В5	A	+0	4.9	3.6	4.0	2.1
	L B						A	+1	5.8 4.6	3.6	5,5 4,3	2.2
							A	+2	5.6	3.7	4.4 5.2	2.5
D2T4N							A	+3	5.6	3.3	5.6 5.4	2.4
3-bit Decoder		12	2.2	9999	D3T8 -0	В5	A B C	_0	2.0	3,9	1.9	3.4
D3T8												

10. OTHERS

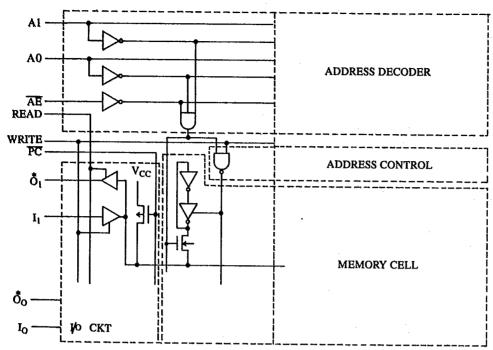
	Macrocell	Equiva-							De	lay		
Function	Fauivalent	lent Gate	Normal- ized	Clamp Level	Symbol	Symbol			t _{PLF}	(ns)	t _{PH}	(ns)
Function Name	Equivalent Circuit	Count	Load Pactor	When Open		No.	Input Name	Output Name	[†] ОLН	K _{LH}	^t OHL	K _{HL}
4-bit Equal Compa- rator ZEQC4		12	1.4	#######	ZEQC4	В5	A ₀ A ₁ A ₂ A ₃ B ₀ B ₁ B ₂ B ₃		11.1	9,0	6.9	1.9
	Pa8						СК		8.8		8.2	
2-bit SR with	CK PRE CK PRE	12	1 1 1,4	@ #:	ZŚRCP3 CK D +A		CLA PRA	+A	13.2	4.4	2.0	1.9
CLR/PRE	D CLR Q D CLR Q QB	12	1 1,4 1	###	PRA CLB +B PRB	В4	CK		8.8		8.2	
ZSRCP3	CLA QA				<u> </u>		CLB PRB	+B	13.2 13.8	4.4	2.0	1.9
2-bit SR	CK CK CK QB	10	1	99	Z\$R -A	B1	СК	+A	4.6	2.4	3.9	1.6
ZSR	D Q Q		,	e (e		D 1	O.R.	+B	5.2	2.2	5,5	1.4

·	Macrocell	Equiva-	Normal-	,		
Function unction Name	Equivalent Circuit	lent Gate Count	ized Load Factor	Symbol	Symbol No.	Delay
Single Port RAM p bits/word v=2.2 word p=2, 4,, 16 v=1, 2,, 32	Ae (a	1.5 · b · w + 3b + 6w + 70		RAMS Aq + Oq	A	60 ns
RAMS				when b-1≥10, the expressions are like follows. b-1		•
Dual Port RAM b bits/words w=2 [®] words b=2, 4,, 16 w=1, 2,, 32	Aa OAa — IAa	2·b·w + 7.5b + 8w + 70		RAMD	A	60 ns
RAM Pre-charge	tpd tpd = 80ns = 20ns	28	1	PCC AE PC	B1	

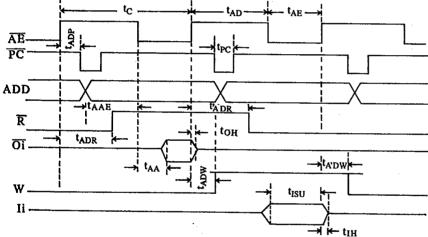
PCC

■ EQUIVALENT CIRCUIT OF RAM (SINGLE PORT RAM)

The following figure is not exactly the same as the actual circuit.







	MIN	TYP	MAX
t _c	400		_
t _{AD}	250	_	T.—
t _{ae}	150		_
t _{ADP}	30	80	160
t _{PC}	10	30	60
t _{AAE}	30		
t _{ADR}	0	_	t _{AD}
t'ADR	0	_	t _{AD}
t _{AA}	20	60	130
t _{on}	4	10	20
t _{abw}	0		t _{AD}
t' _{ADW}	0	_	t _{AD}
t _{isu}	30	_	-
t _{IH}	30	_	

FUNCTIONAL TEST FOR ON-CHIP RAM

In order to easily test on-chip RAM, the logic design should be done so as to access the RAM directly outside the chip, for instance in RAM-TEST MODE.

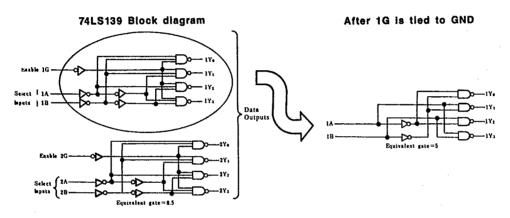
■ TTL TO CMOS GATE ARRAY CONVERSION

(1) Equivalent Gate Count

No.	TTL	Gate Count	No.	TTL	Gate Count
1	74LS00	1	41	74LS90	42
2	74LS01	1	42	74LS91	66
3	74LS02	ī	43	74LS92	50.5
4	74LS03	1	44	74LS93	50.5
5	74LS04	0.5	45	74LS95	53
٠, ٢	74LS05	0.5	46	74LS107	13
6	74LS08	1.5	47	74LS109	13
7	74LS09	1.5	48	74LS112	13.5
8	74LS10	1,5	49	74LS113	13
9	74LS11	2	50	74LS114	13.5
10	74LS12	1.5	51	74LS125	2
11	74LS15	2	52	74LS126	2
12	74LS20	2	53	74LS136	2.5
13	74LS21	2.5	54	74LS137	31
14	74LS22	2	55	74LS138	19.5
15	74LS26	1	56	74LS139	7,5
16	74LS27	1.5	57	74L\$145	22
17	74LS28	1	58	74LS147	36
18	74LS30	5,5	59	74LS148	38
19	74LS32	1.5	60	74LS151	45.5
20	74LS33	1	61	74LS152	25
21	74LS37	1	62	74LS153	23
22	74LS38	1	63	74LS154	77
23	74LS40	2	64	74LS155	16.5
24	74LS42	24	65	74LS156	16.5
25	74LS47	40	66	74LS157	12.5
26	74LS48	42	67	74LS158	10.5
27	74LS49	35	68	74LS160	64.5
28	74LS51	3.5	69	74LS161	64.5
29	74LS54	6.5	70	74LS162	58
30	74LS55	4,5	71	74LS163	64
31	74LS63	1	72	74LS164	82
32	74LS73	13	73	74LS166	102
33	74LS74	9	74	74LS170	114
34	74LS75	7	75	74LS174	49.5
35	74LS76	13,5	76	74LS175	33.5
36	74LS77	7	77	74LS181	93
37	74LS78	13.5	78	74LS183	12
38	74LS83	52,5	79	74LS190	84.5
39	74LS85	63	80	74LS191	78.5
40	74LS86	2.5	81	74LS192	63

No.	TTL	Gate Count	No.	TTL	Gate Coun
82	74LS193	63	97	74LS283	51,5
83	74LS194	77,5	98	74LS290	41.5
84	74LS195	56	99	74LS298	43.5
00	74LS246	40.5	100	74LS299	144
85	74LS247	48.5	101	74LS365	25
86	74LS248	48.5	102	74LS366	25
87	74LS249	48.5	103	74LS367	12.5
88	74LS251	39	104	74LS368	12.5
89	74LS253	25.5	105	74LS375	7
90	74LS257	21.5	106	74LS386	2.5
91	74LS258	19	107	74LS390	37.5
92	74LS259	79	108	74LS393	33.5
93	74LS266	2,5	109	74LS490	40.5
94	74LS273	65,5	110	74LS668	82
95	74LS279	4	111	74LS669	71.5
96	74LS280	45,5	112	74LS670	115.5

(2) Figures indicated in the table are the guidelines when all pin functions are fully converted.



Note: Exact conversion is not necessarily achieved since certain pins are tied to GND level, such as Enable 1G pin of LS139 shown above. In this case, the equivalent gate count will be off by 2.5 gates.

Therefore, users might overestimate the gate count by 20 to 30% if the figures above are employed without due consideration.

HIT I HOLL AND THE TO

T-42-11-09

■ CMOS LOGIC TO CMOS GATE ARRAY

(1) Equivalent Gate Count

Gate Count CMOS Gate Count No. **CMOS** No. 41 HD14161B HD14001B 68 1 2 42 HD14162B 67 2 HD14002B 43 68 3 HD14006B 144.5 HD14163B 44.5 44 HD14174B 48.5 4 HD14008B 5 HD14011B 1 45 HD14175B 32.5 46 HD14194B 65.5 HD14012B 2 6 7 8 47 HD14501UB 2, 1.5 HD14013B 69.5 48 HD14502B 18.5 8 HD14014B 9 32.5 49 HD14503B 13 HD14015B 50 HD14506B 11 10 N/A HD14017B 116 51 HD14508B 17.5 11 HD14020B 52 HD14510B 78.5 50.5 12 HD14021B 20.5 13 HD14023B 1.5 53 HD14512B 55 54 HD14514B 57 14 HD14024B 1.5 55 HD14515B 63 15 HD14025B 56 HD14516B 78.5 16 HD14027B 12 57 HD14517B 563.5 17 HD14028B 26 39 58 HD14518B 18 HD14032B 68.5 59 HD14519B 26 19 HD14035B 57.5 38 HD14520B 80 60 20 HD14038B N/A 61 HD14529B 21 21 HD14040B 30 22 HD14042B 16.5 62 HD14531B HD14532B 41.5 23 HD14043B 20.5 63 N/A 20.5 64 HD14538B 24 HD14044B 24 0.5 65 HD14539B 25 HD14049UB N/A 26 HD14050B 1 66 HD14541B 7.5 5.5 67 HD14555B 27 HD14068B 8.5 28 HD14069UB 0.5 68 HD14556B 53 29 HD14070B 2.5 69 HD14558B 68 70 HD14560B 30 HD14071B 1.5 71 768.5 2.5 HD14562B 31 HD14072B 1, 0.5 2 72 HD14572UB 32 HD14073B 12 2 73 HD14583B HD14075B 33 74 N/A 34 HD14076B 53 HD14584B 2,5 75 HD14585B 37.5 35 HD14077B HD14078B 5.5 36 HD14081B 1.5 37 HD14082B 2.5 38 N/A 39 HD14093B

67

40

HD14160B

■ PACKAGE OUTLINE

